

Data-Oriented Programming

On the Expressiveness of Non-Control Data Attacks

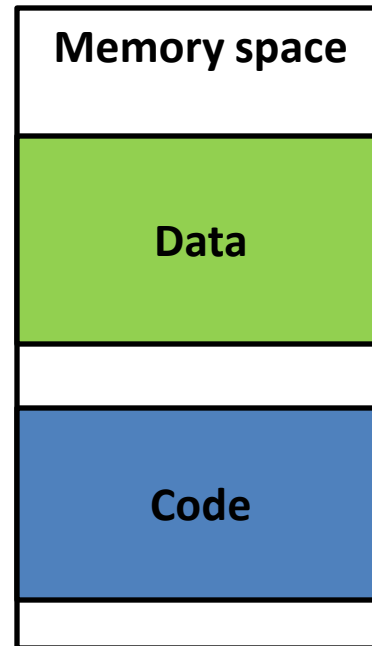
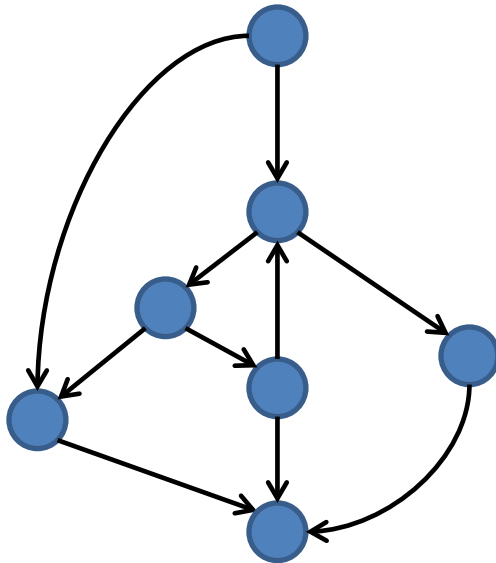
Hong Hu, Shweta Shinde,
Sendroiu Adrian, Zheng Leong Chua,
Prateek Saxena, Zhenkai Liang

*Department of Computer Science
National University of Singapore*

Control Attacks are Getting Harder

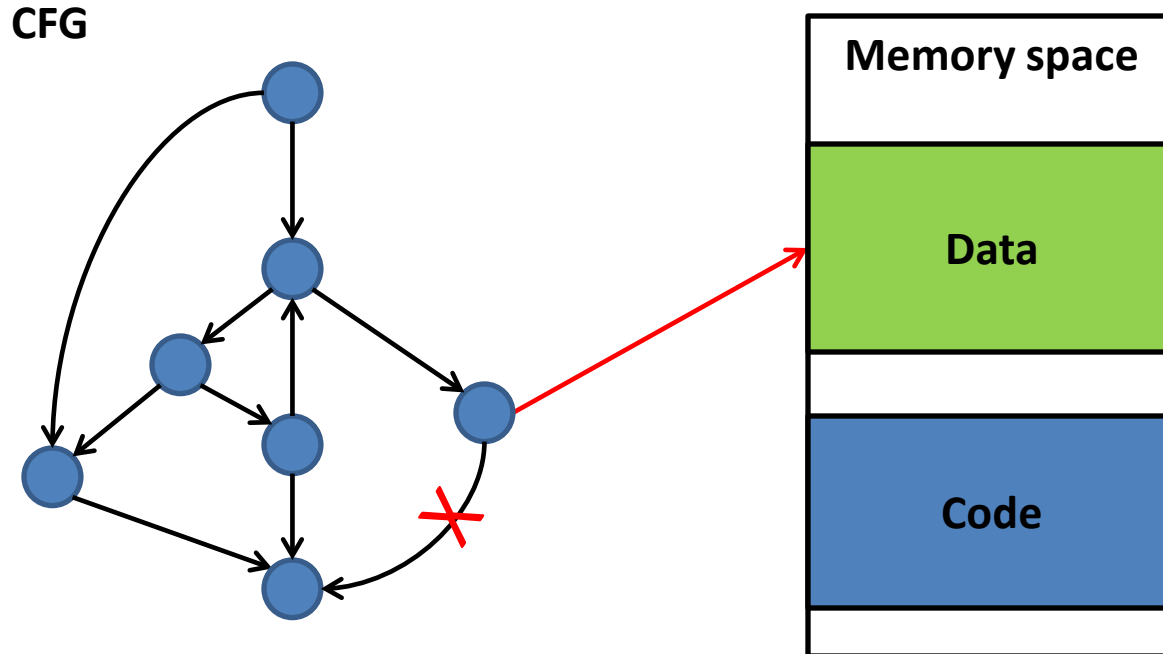
Control Attacks are Getting Harder

CFG



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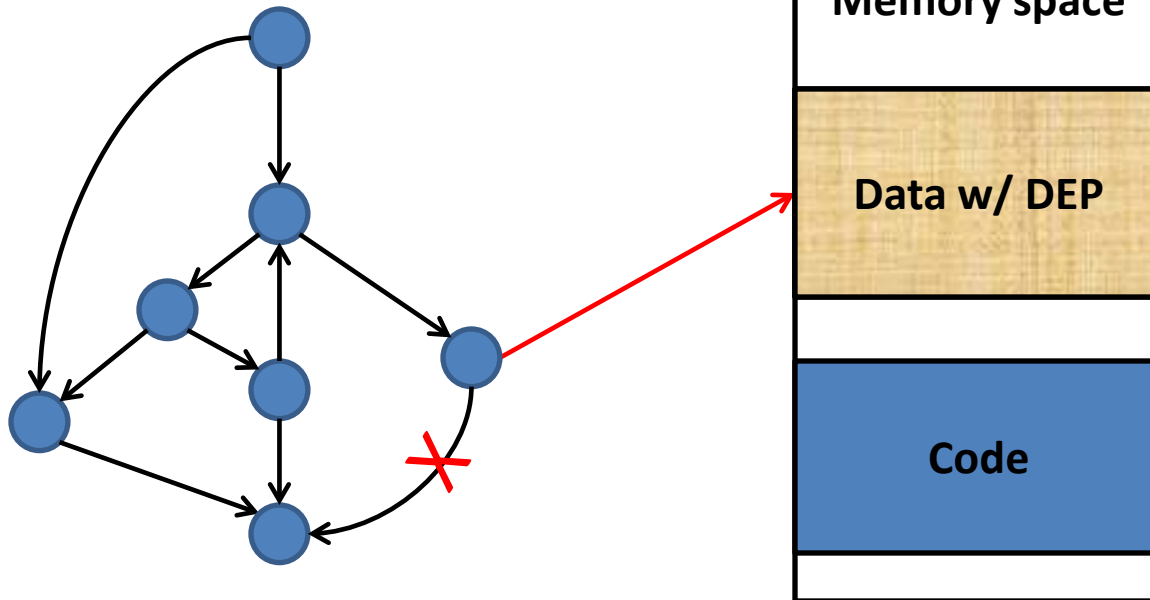
- Code injection



Control Attacks are Getting Harder

- Code injection ← Data Execution Prevention

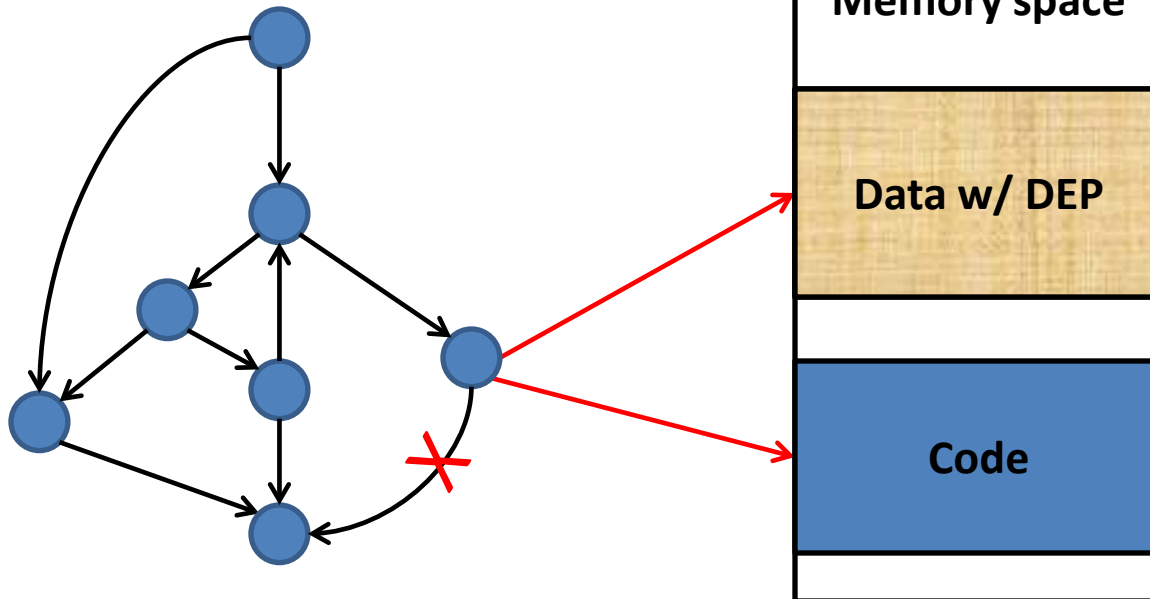
CFG



Control Attacks are Getting Harder

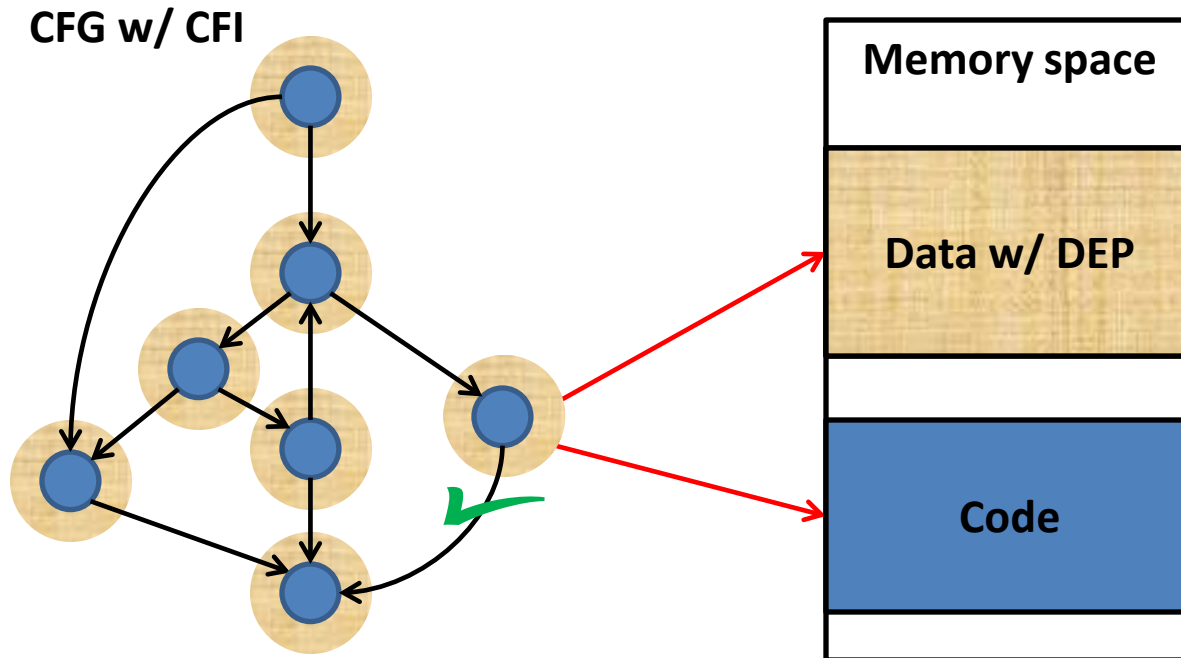
- Code injection ← Data Execution Prevention
- Code reuse
 - return-to-libc
 - return-oriented programming (ROP)

CFG



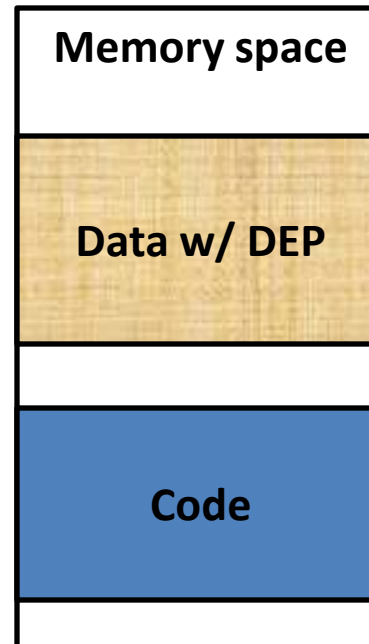
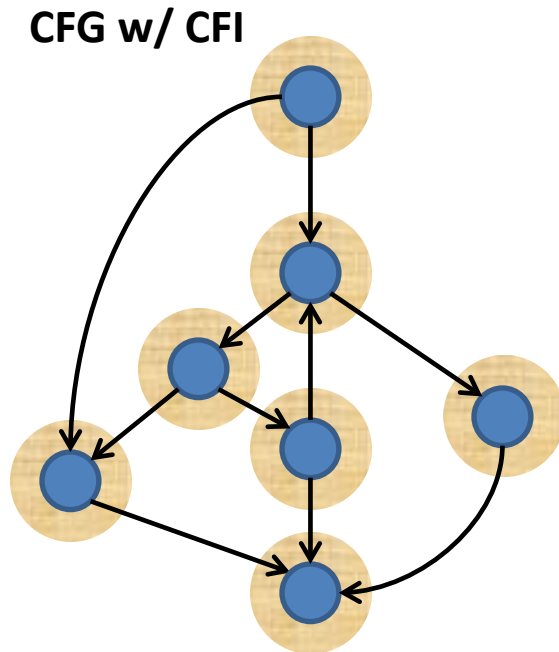
Control Attacks are Getting Harder

- Code injection ← Data Execution Prevention
- Code reuse ← Control Flow Integrity
 - return-to-libc
 - return-oriented programming (ROP)



A New Attack Class

- Assume: conform to CFI & DEP

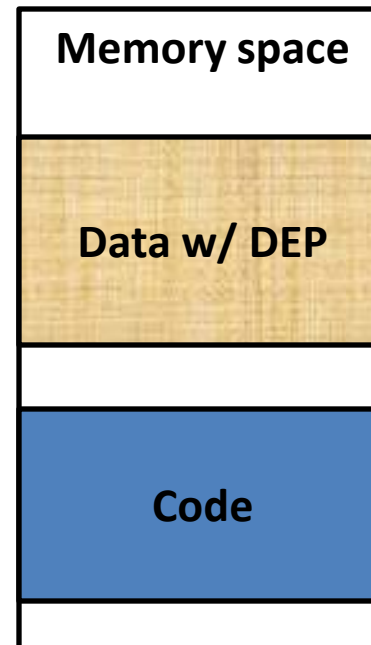
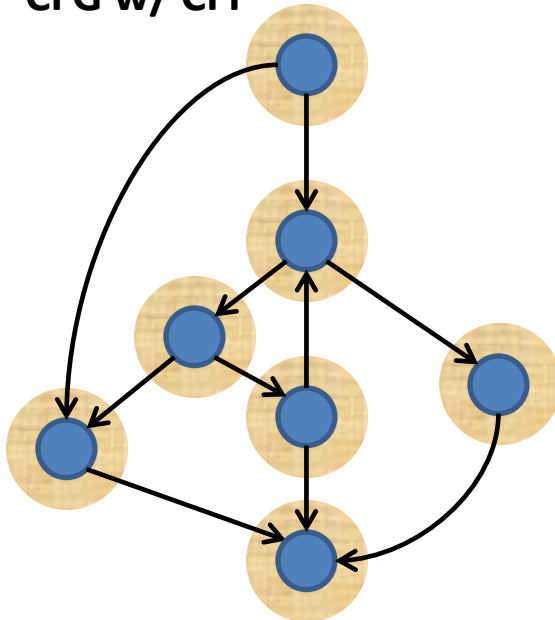


A New Attack Class

- Assume: conform to CFI & DEP
- Attackers' capability on arbitrary vul. programs?



CFG w/ CFI



Non-Control Data Attacks

- Corrupt/leak several bytes of security-critical data

Non-Control Data Attacks

- Corrupt/leak several bytes of **security-critical data**

```
//set root privilege  *
seteuid(0);
.....
//set normal user privilege
seteuid(pw->pw_uid);
//execute user's command
```

```
//offset depends on IE version +
safemode = *(DWORD *)
            (jsobj + offset);
if(safemode & 0xB == 0) {
    Turn_on_God_Mode();
}
```

* Shuo Chen, Jun Xu, Emre C. Sezer, Prachi Gauriar, and Ravishankar K. Iyer. Non-Control-Data Attacks Are Realistic Threats. In USENIX 2005.

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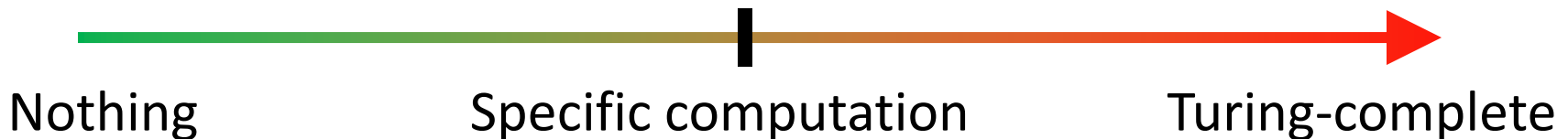
Non-Control Data Attacks

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- Special cases relying on particular data/functions
 - user id, safemode, private key, etc
 - interpreter – printf() (with “%n”), etc



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- Data-Oriented Programming (DOP)
 - build expressive non-control data attacks
 - independent of any specific data / functions
- DOP builds attacks on real-world programs
 - bypass ASLR w/o address leakage
 - simulate a network bot
 - enable code injection

Motivating Example

```
1 struct server{int *cur_max, total, typ;} *srv;
2 int quota = MAXCONN; int *size, *type;
3 char buf[MAXLEN];
4 size = &buf[8]; type = &buf[12]
5 ...
6 while (quota-- ) {
7     readData(sockfd, buf);          // stack bof
8     if(*type == NONE ) break;
9     if(*type == STREAM)
10         *size = *(srv->cur_max);
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Vulnerable
Program

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Vulnerable
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1 struct Obj {struct Obj *next; int prop;}
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3 void updateList(struct Obj *list, int addend){
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Malicious
Computation

Motivating Example

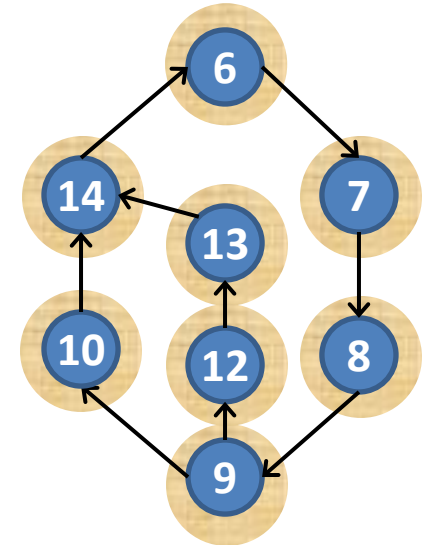
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CFG w/ CFI



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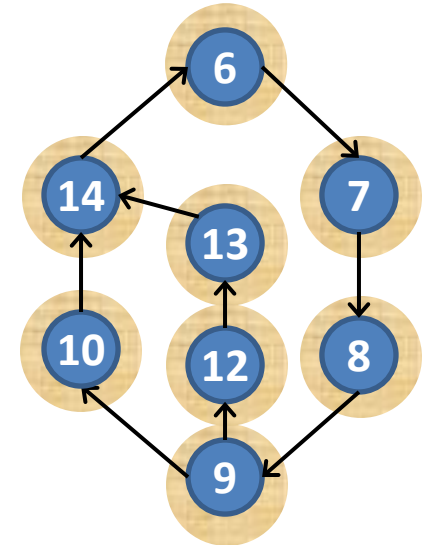
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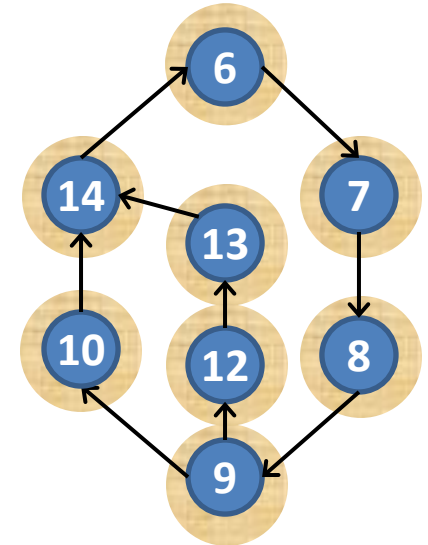
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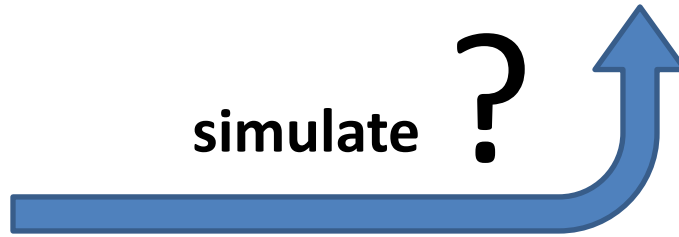
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malicious computation

simulate ?



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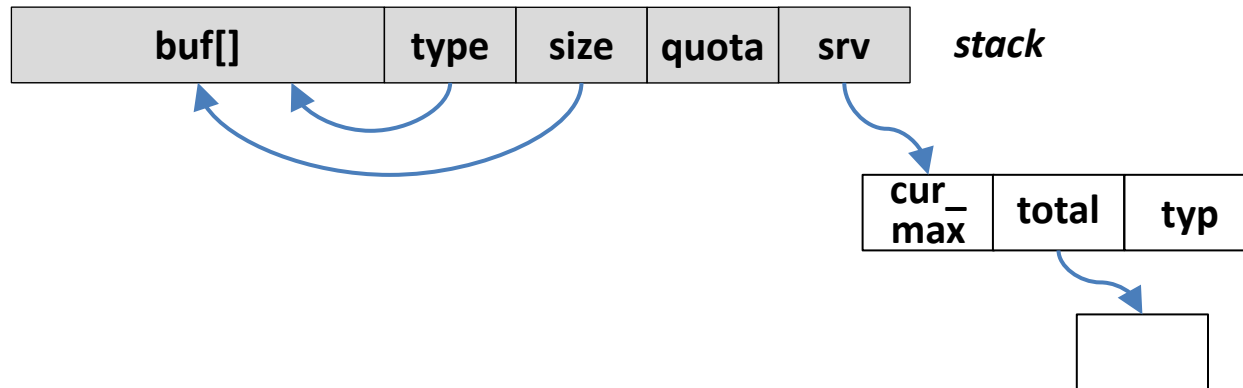
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Memory space



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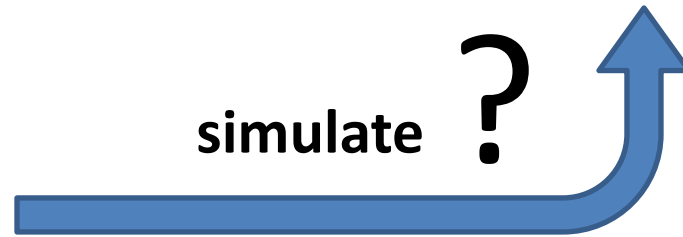
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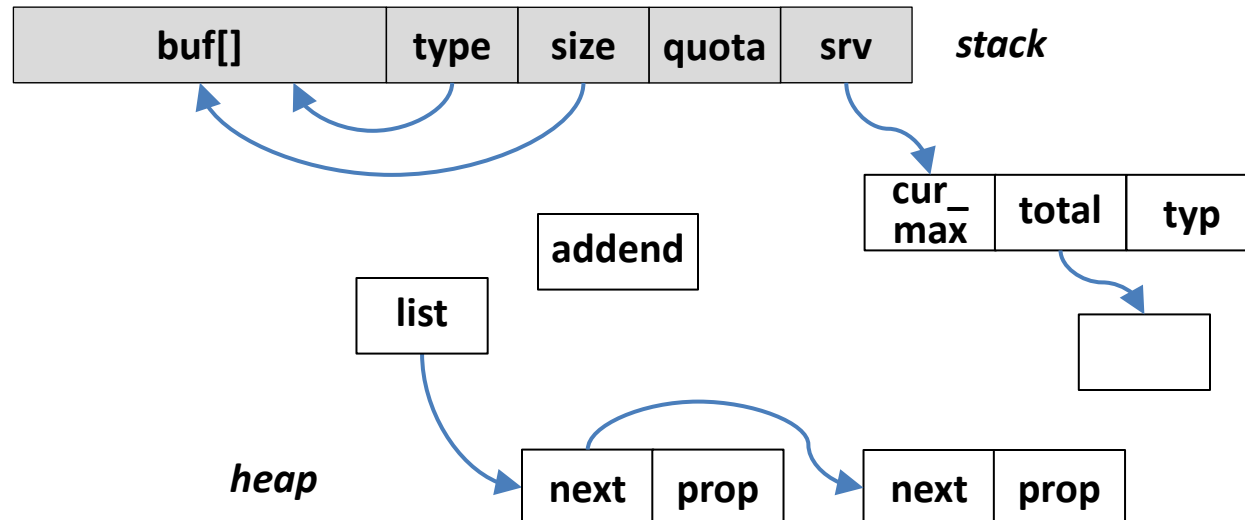
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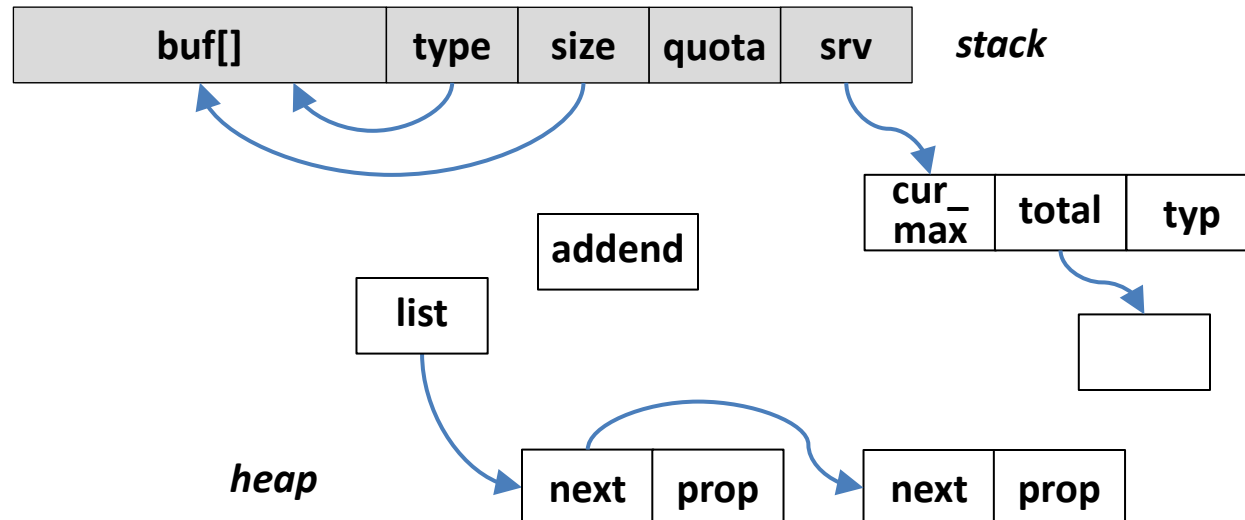
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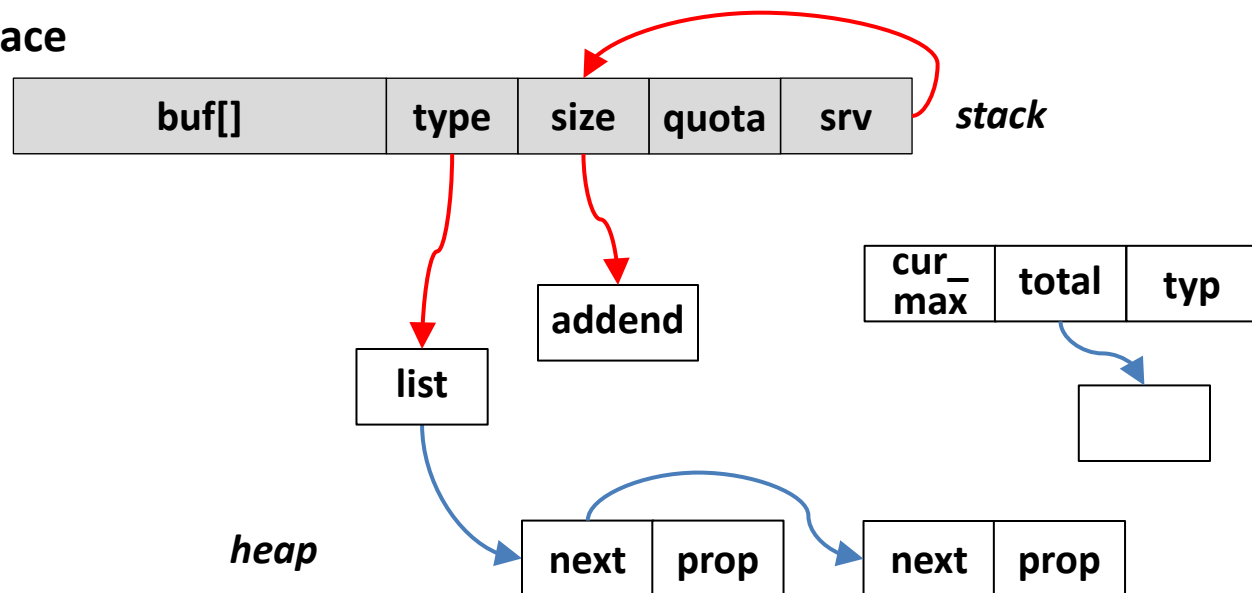
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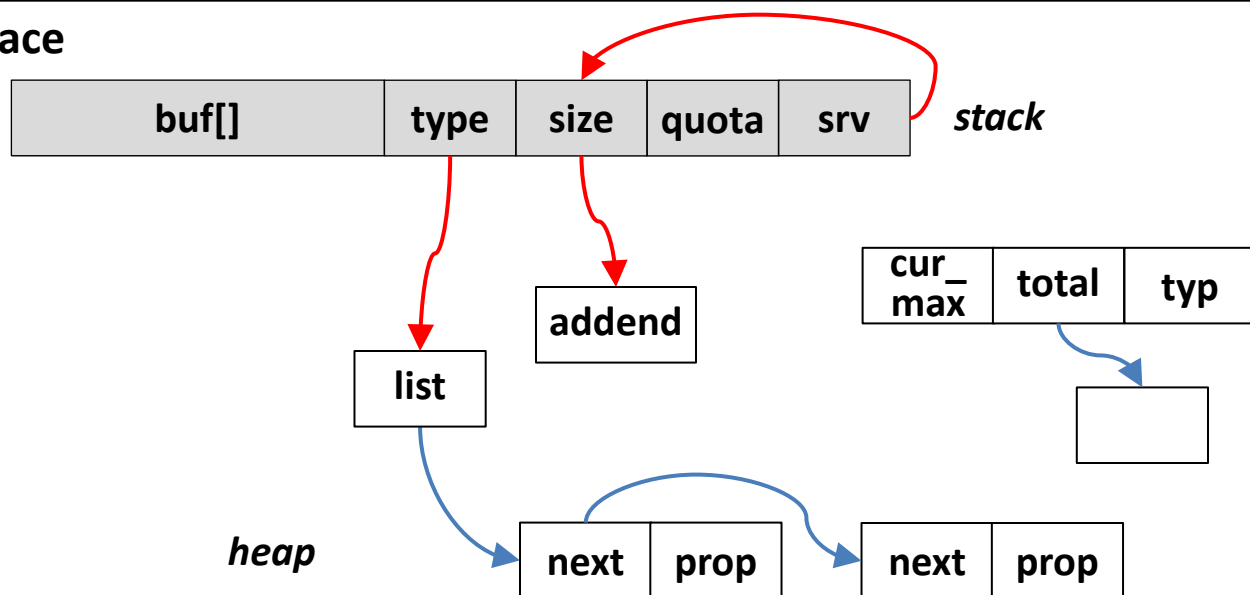
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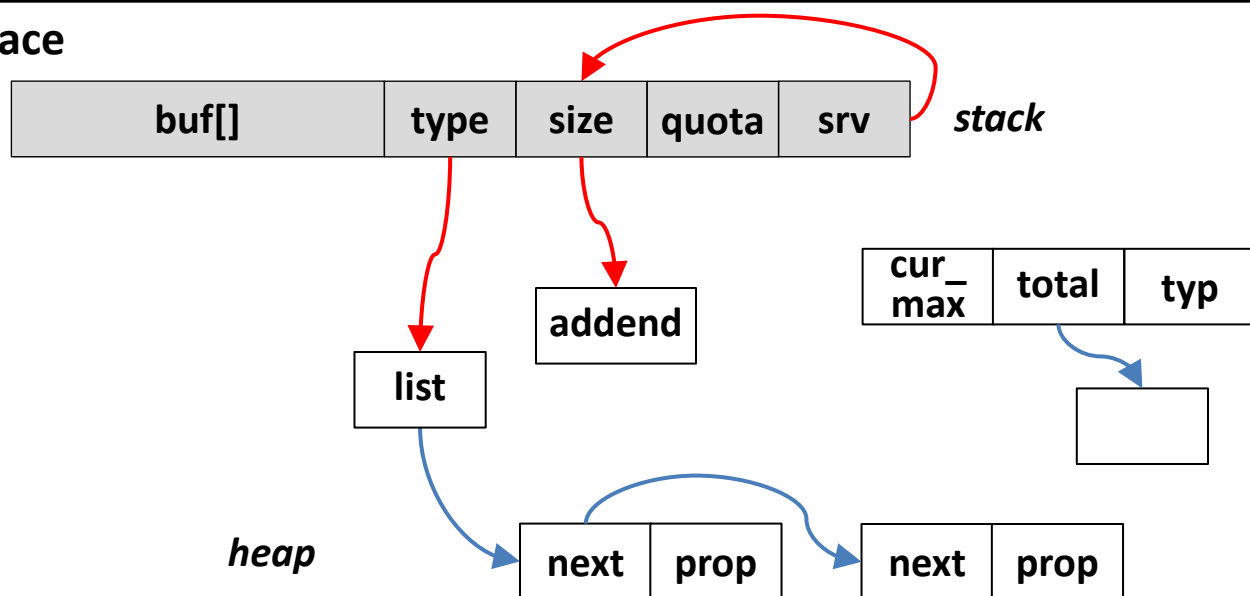
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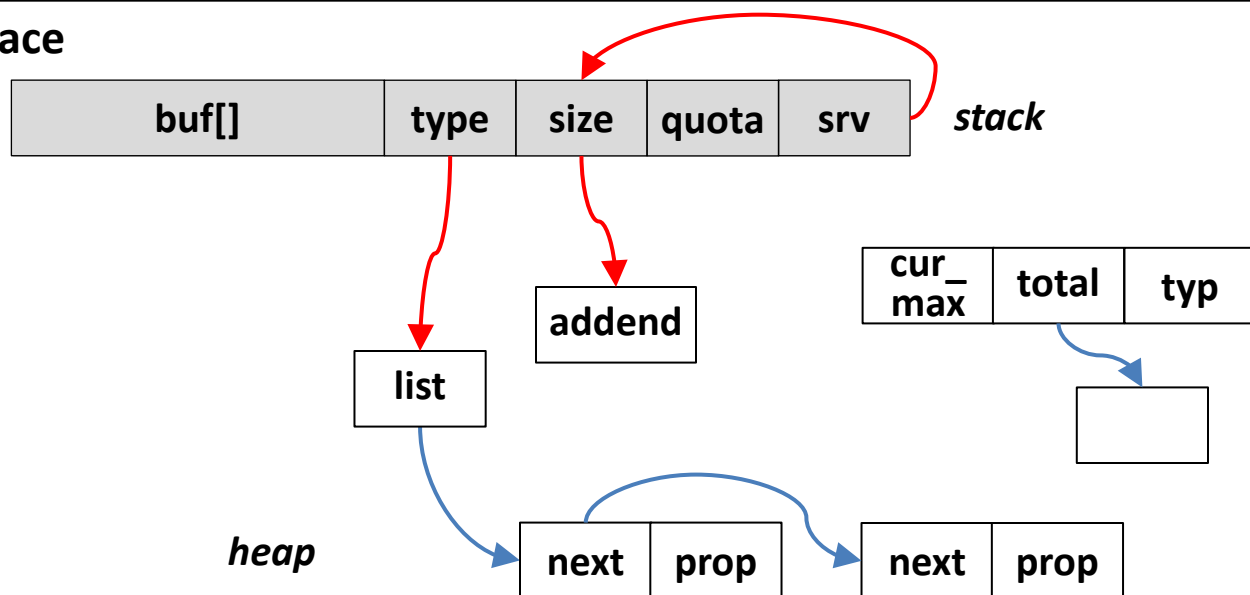
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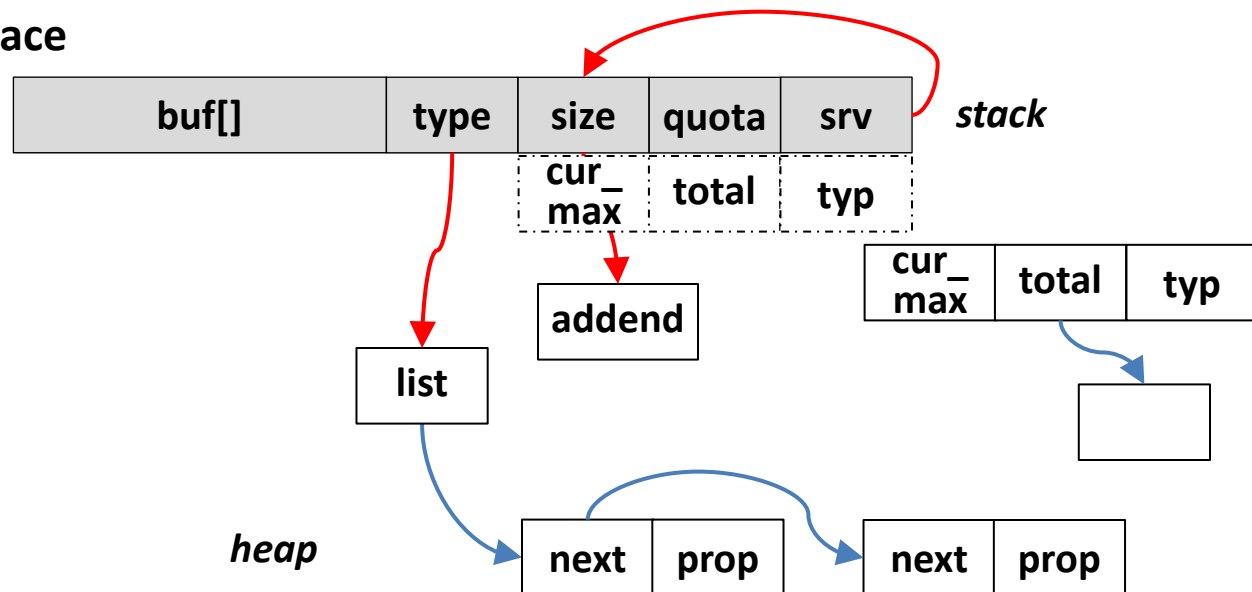
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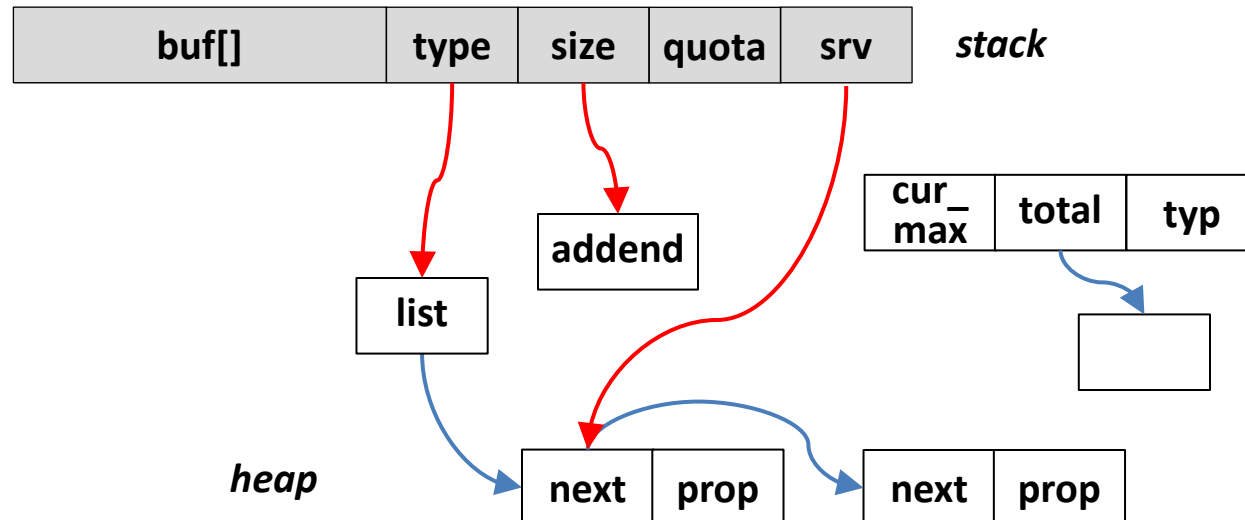
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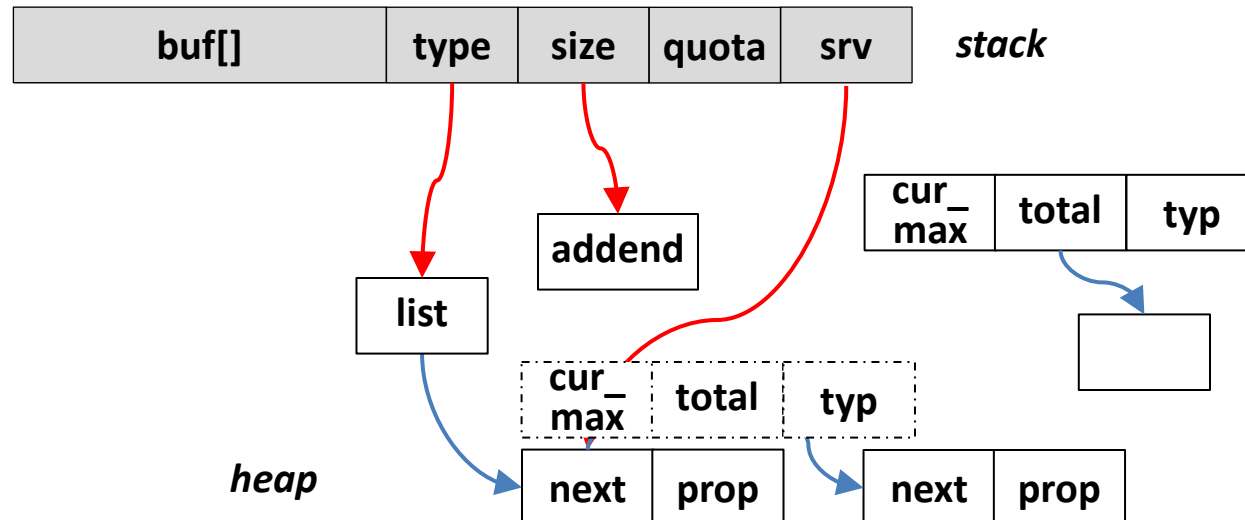
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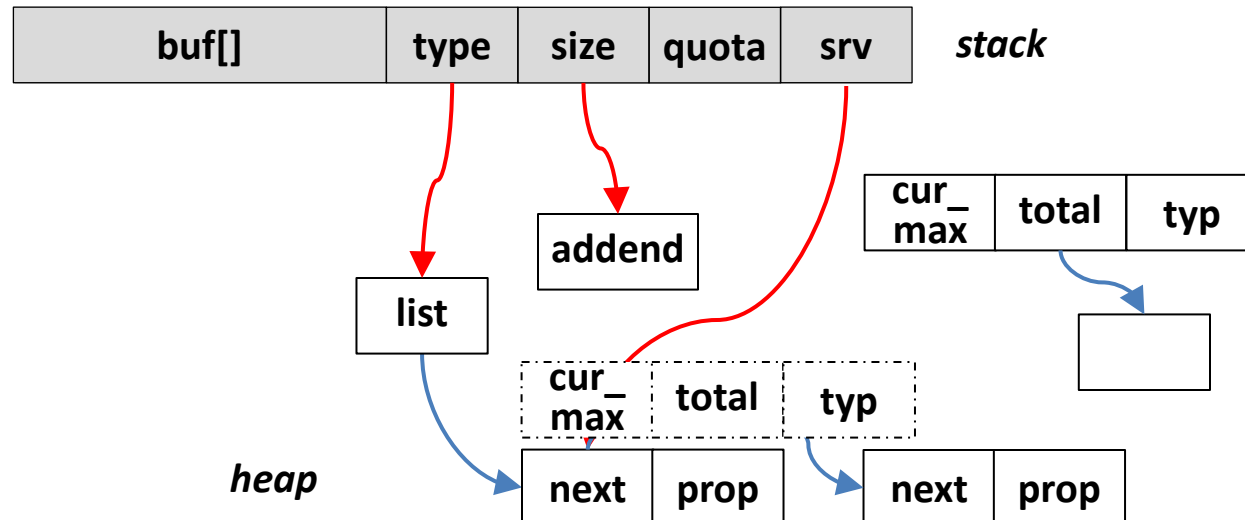
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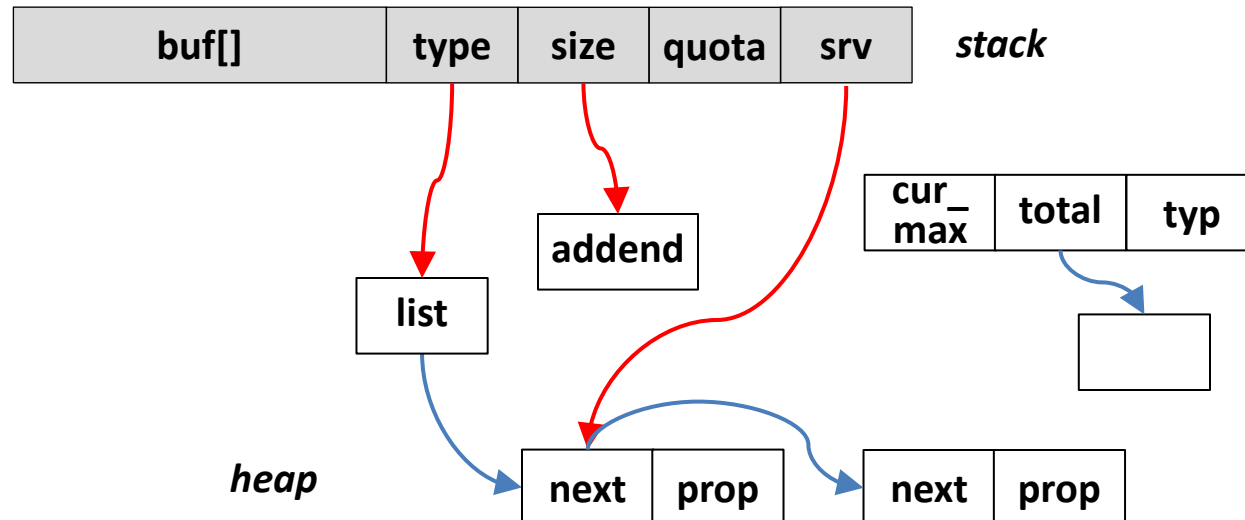
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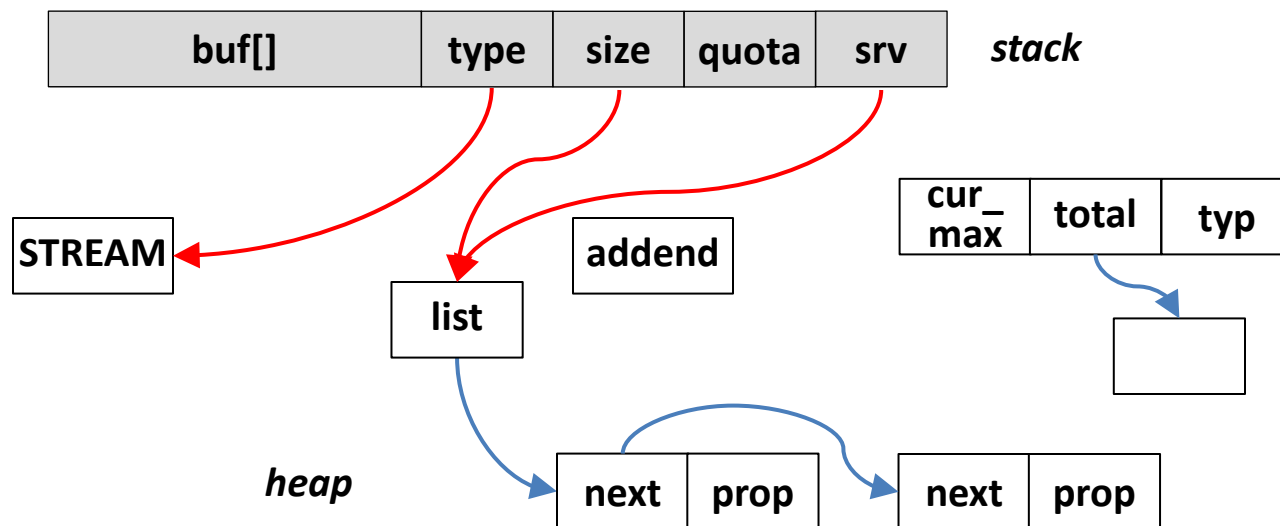
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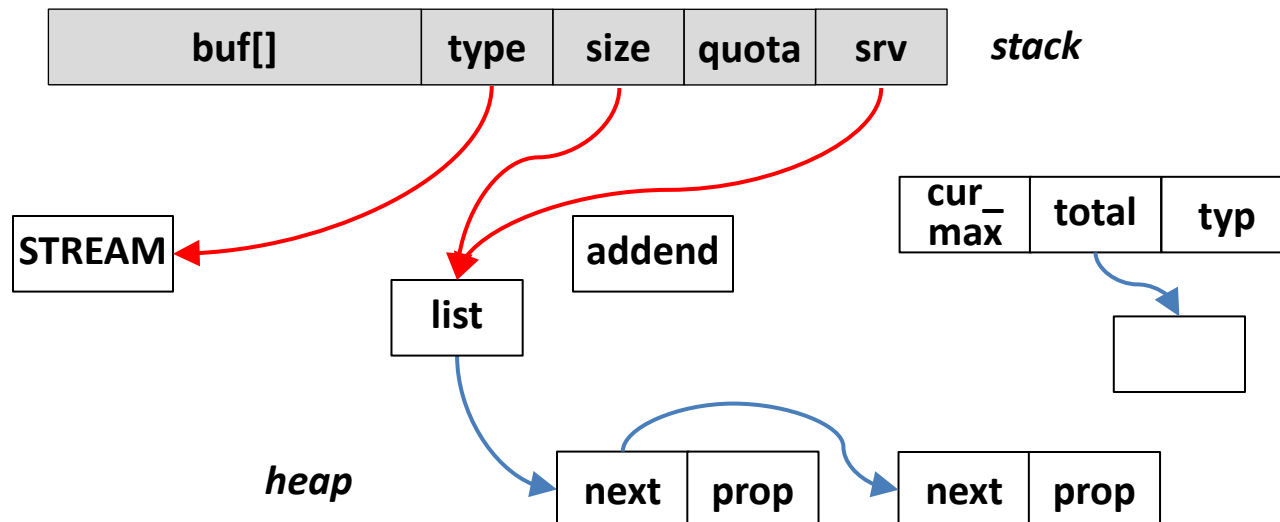
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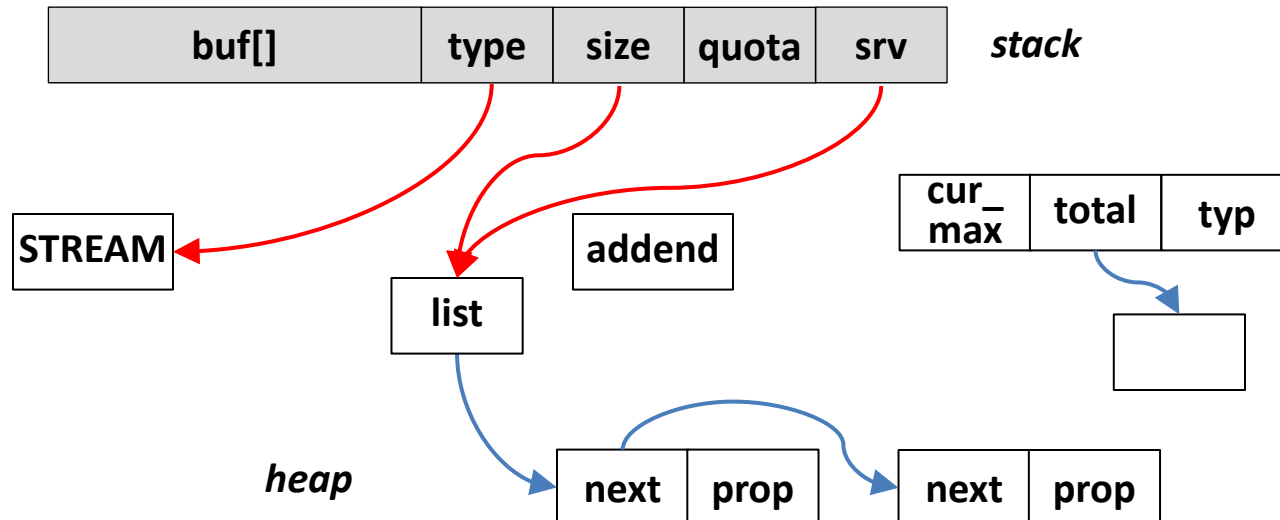
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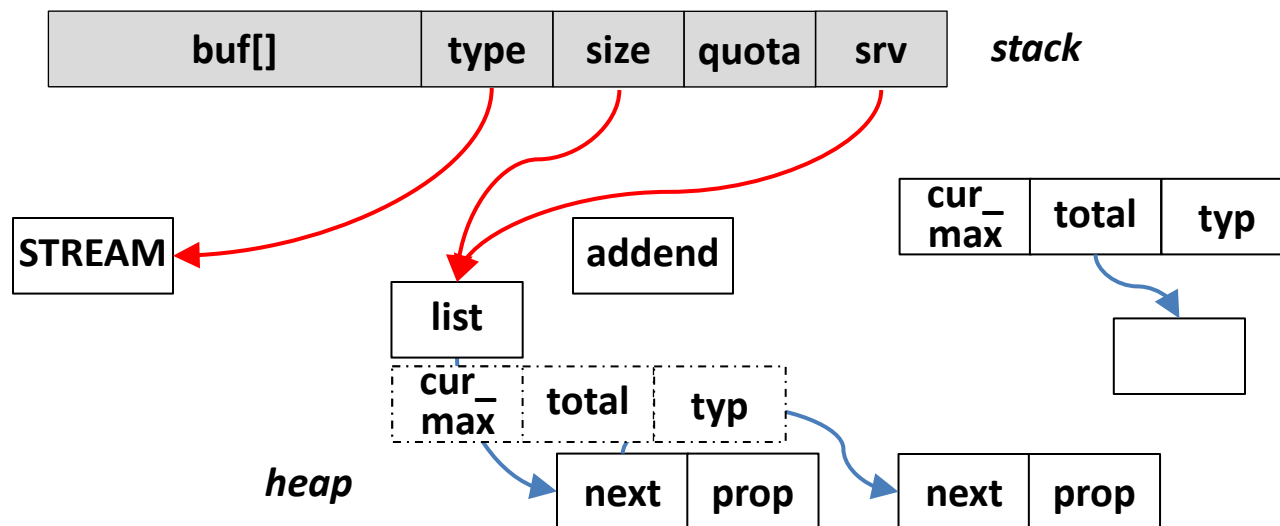
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5   list->prop += addend;
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malicious computation

simulate ?

Memory space



Motivating Example (cont.)

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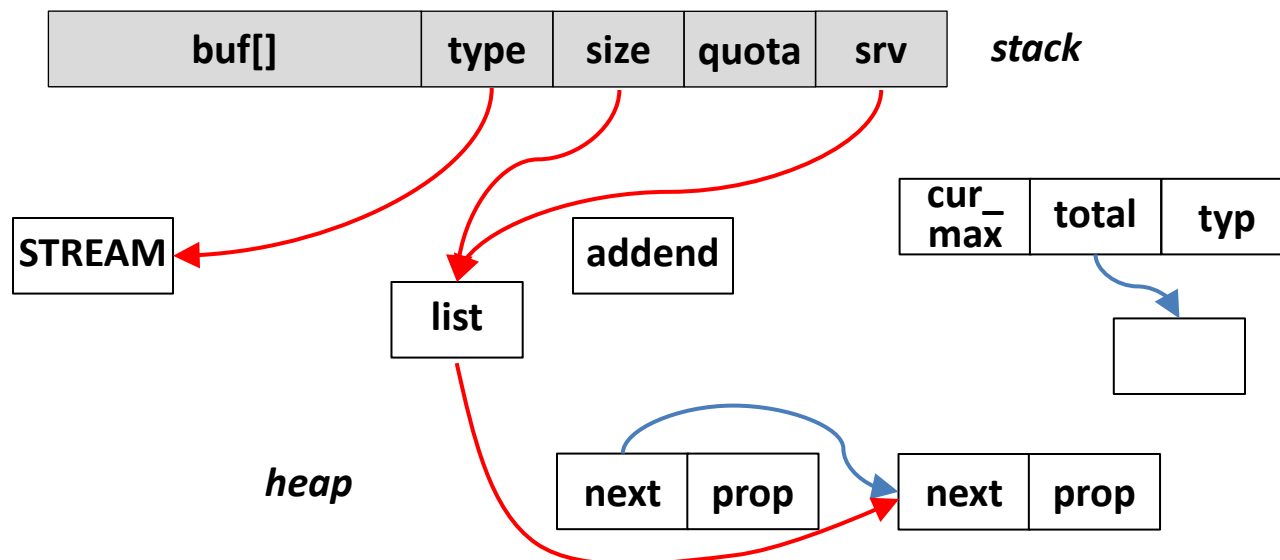
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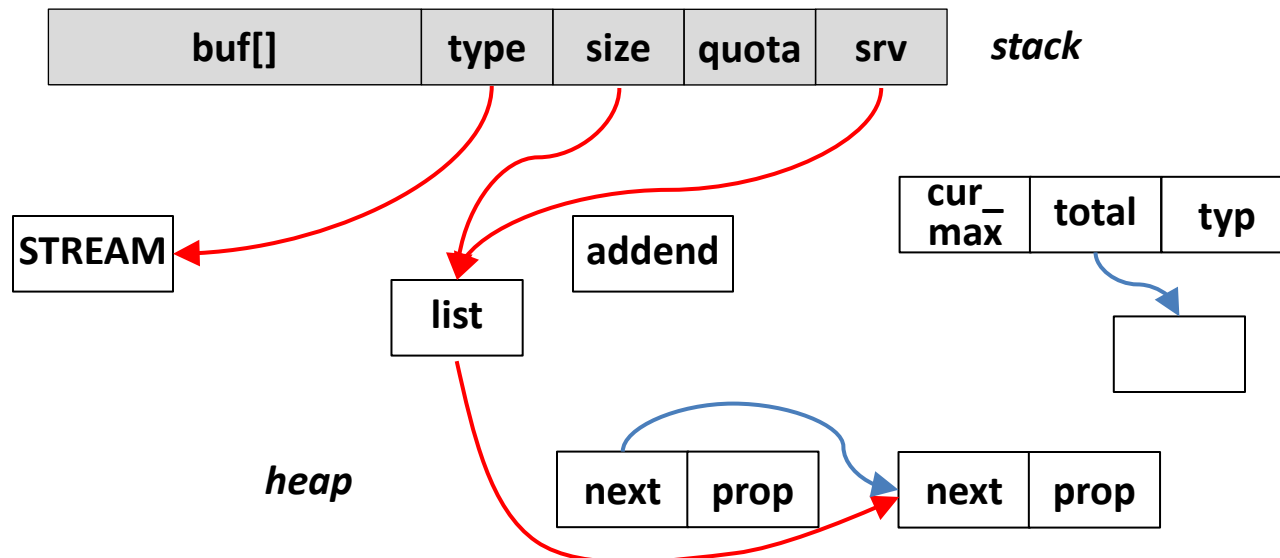
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Data-Oriented Programming

A Generic Technique

Data-Oriented Programming (DOP)

- General construction
 - w/o dependency on specific data / functions

Data-Oriented Programming (DOP)

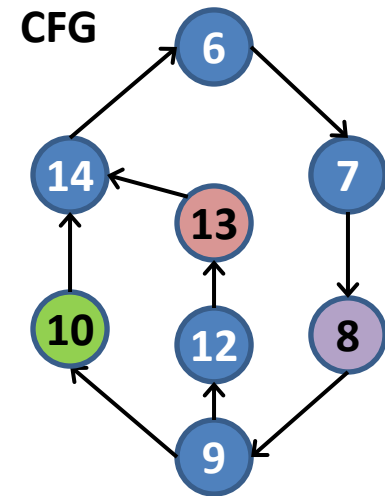
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 - w/o dependency on specific data / functions
- Expressive attacks
 - towards Turing-complete computation

Data-Oriented Programming (DOP)

- General construction
 - w/o dependency on specific data / functions
- Expressive attacks
 - towards Turing-complete computation
- Elements
 - data-oriented gadgets
 - gadget dispatchers

Data-Oriented Gadgets

- x86 instruction sequence
 - show in normal execution (CFI)

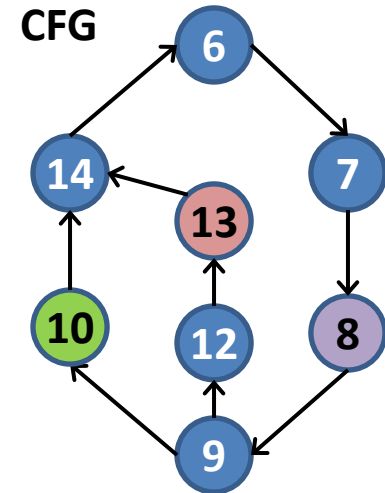


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Addition: `srv->total += *size;`

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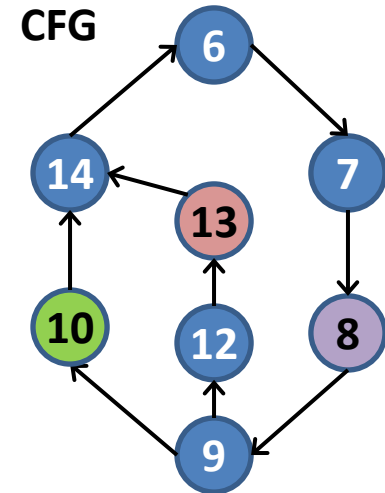


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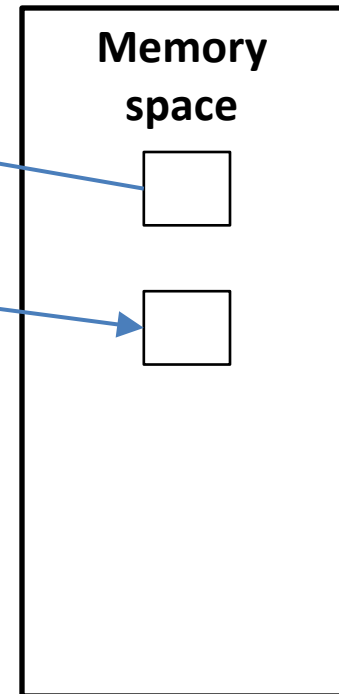
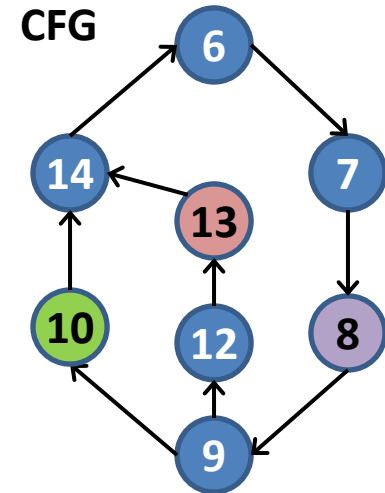


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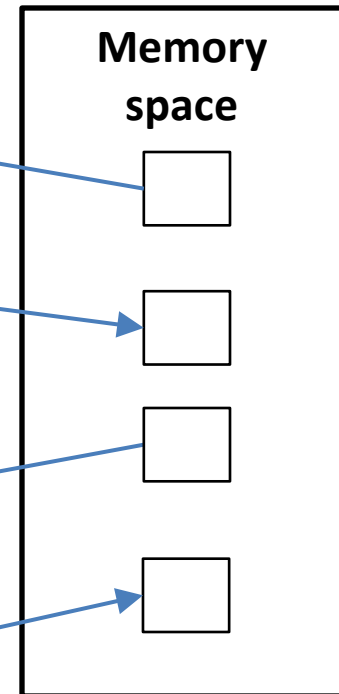
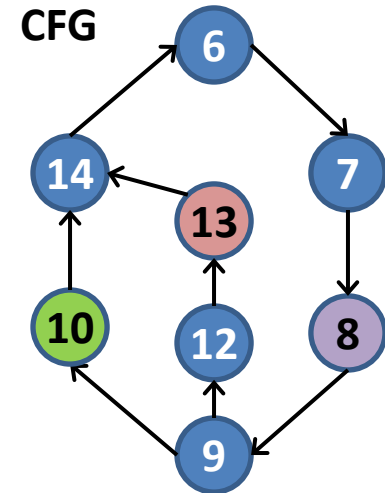
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Data-Oriented Gadgets

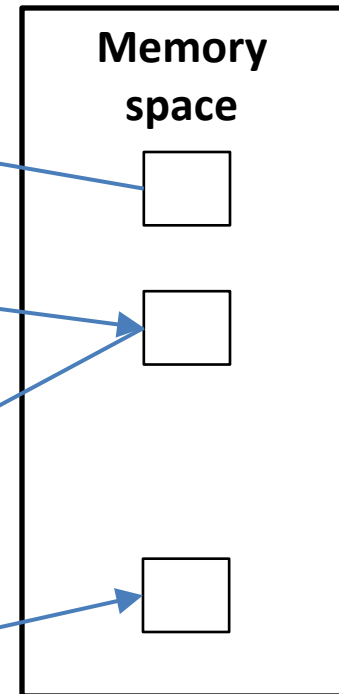
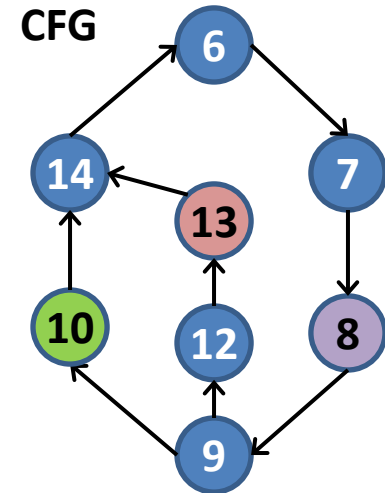
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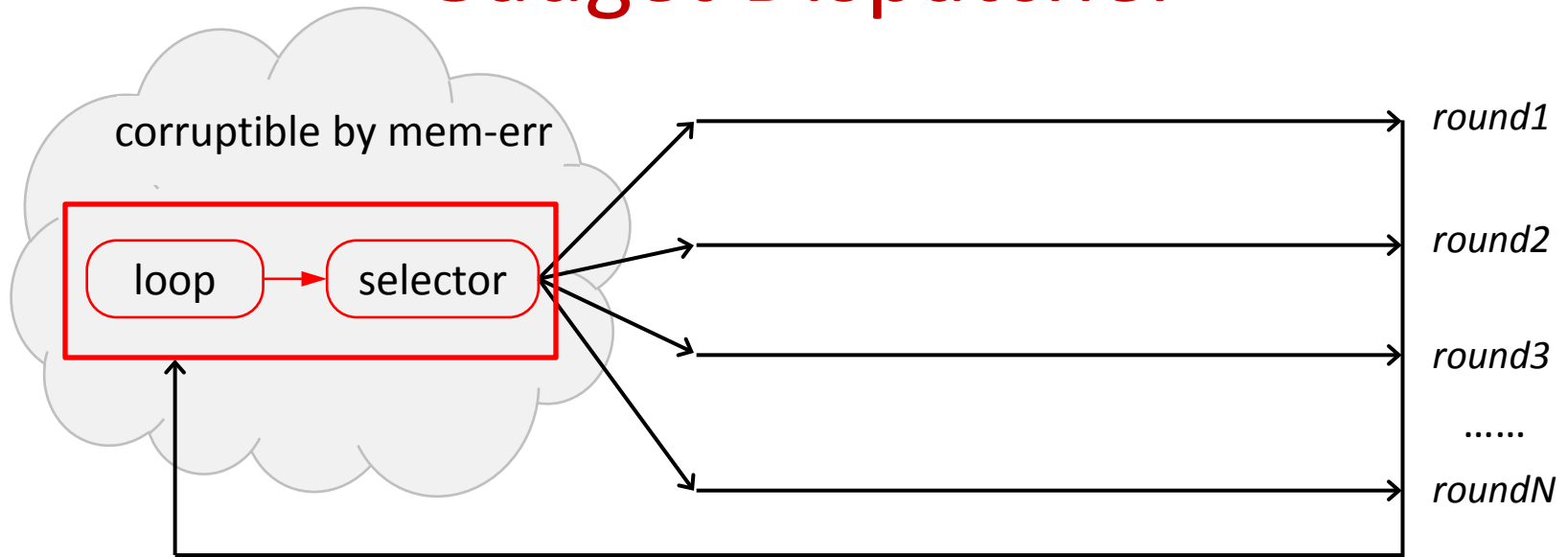
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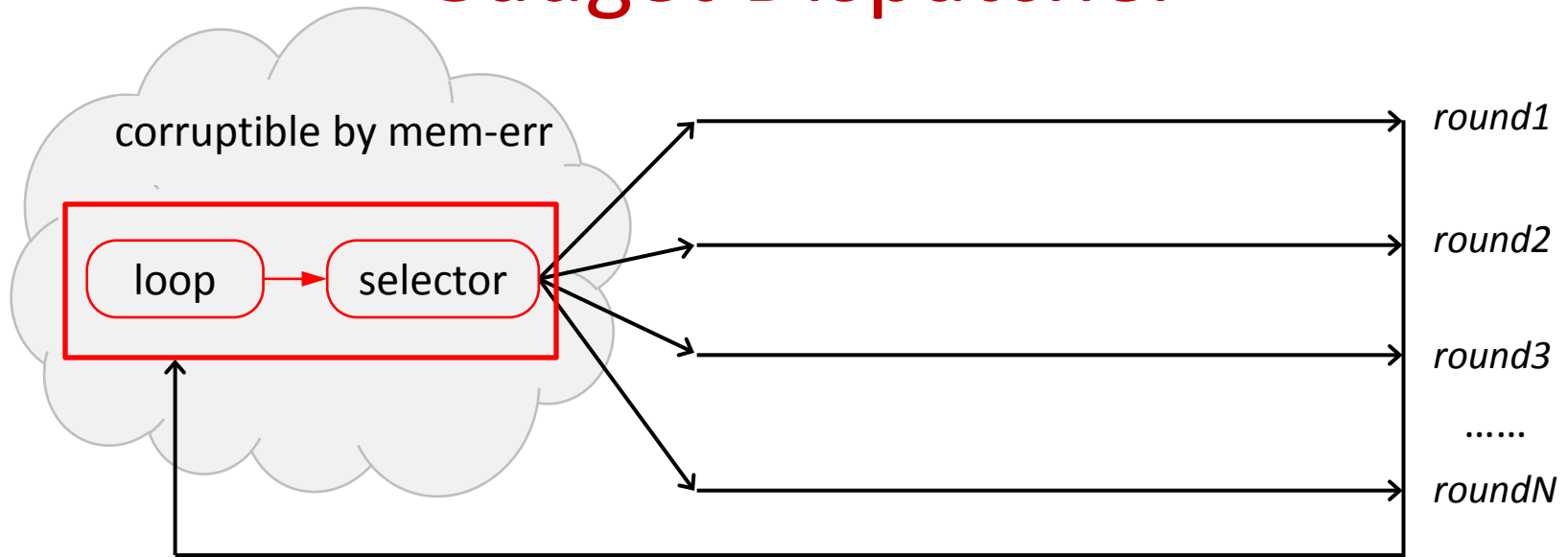
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Gadget Dispatcher

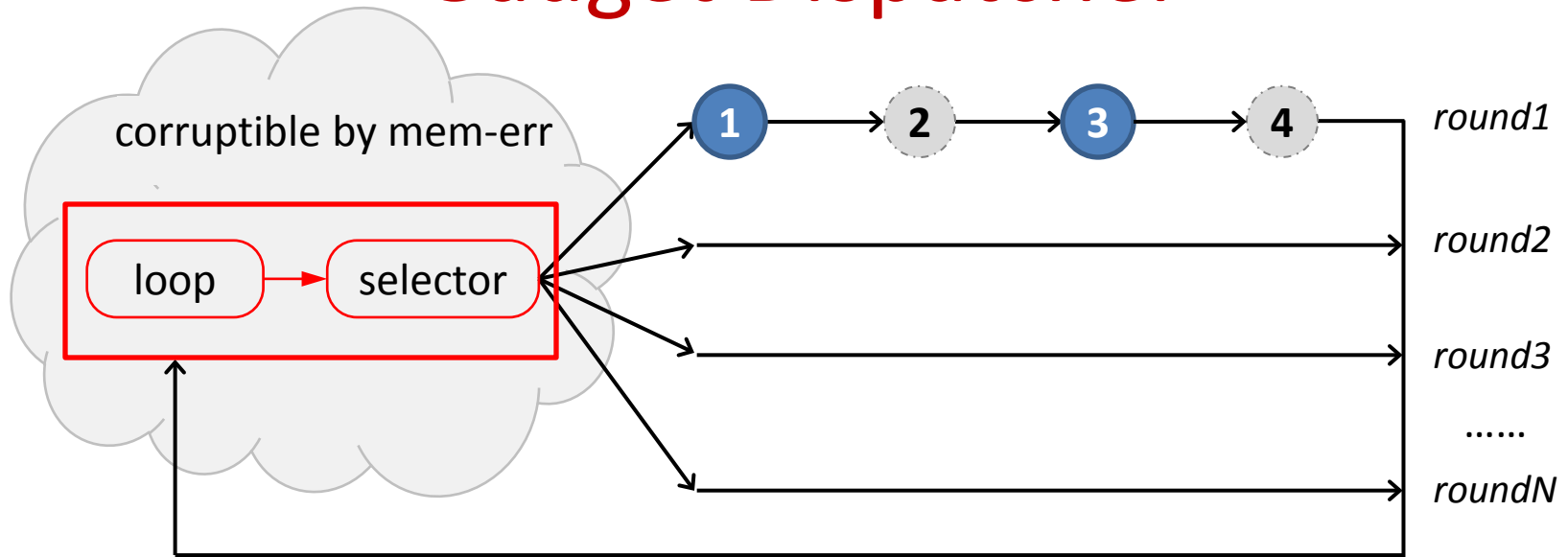


Gadget Dispatcher



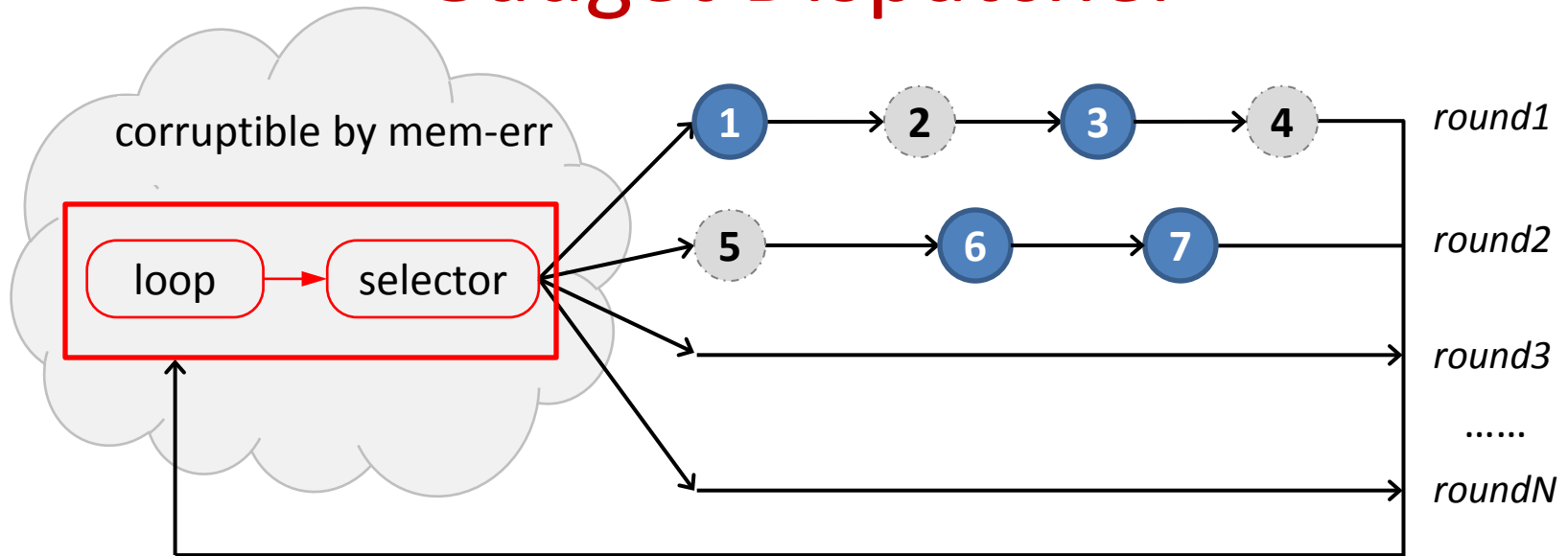
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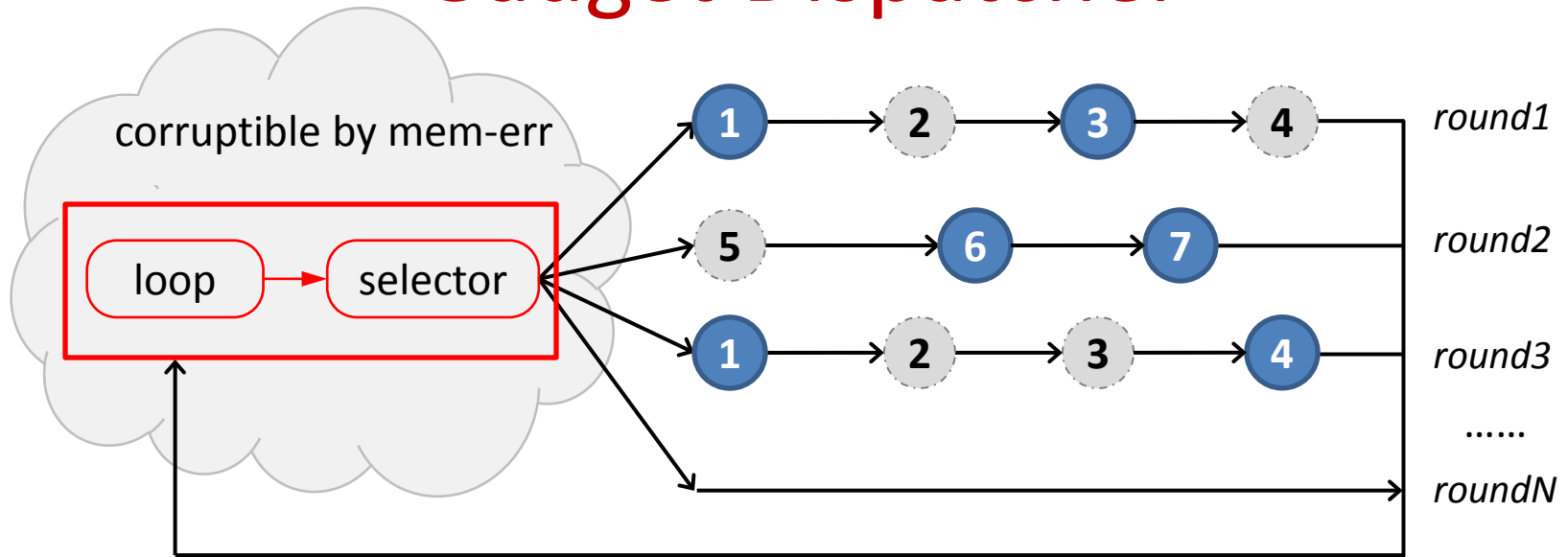
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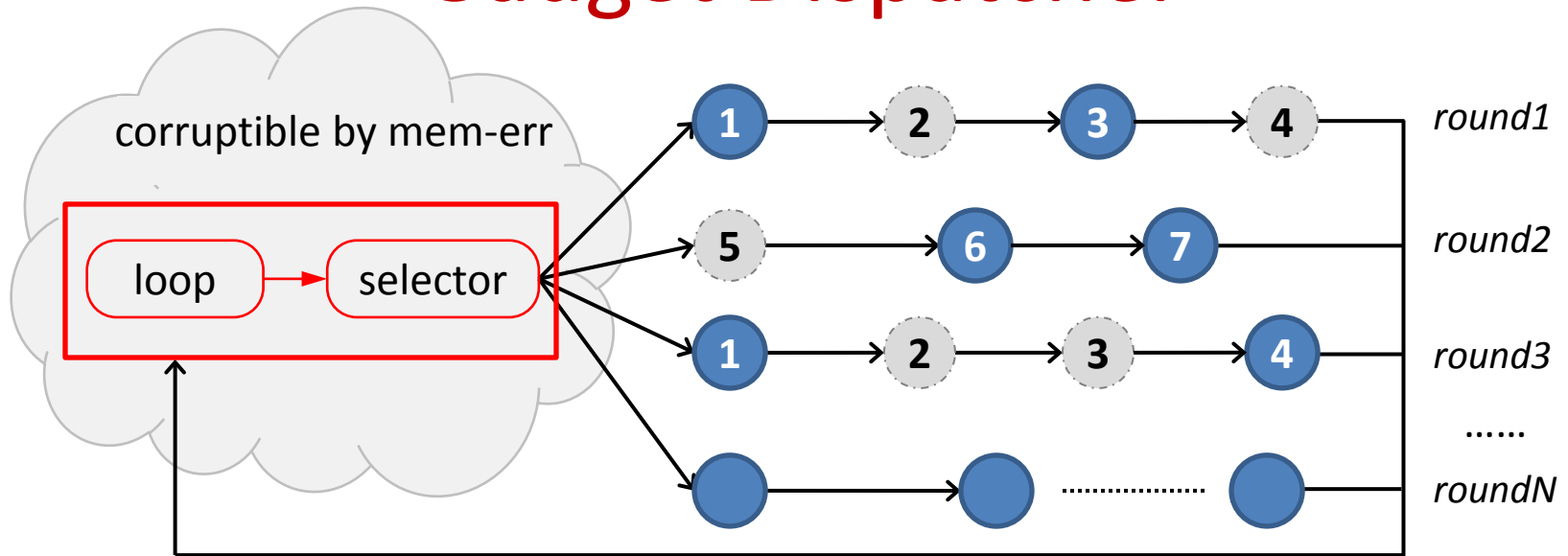
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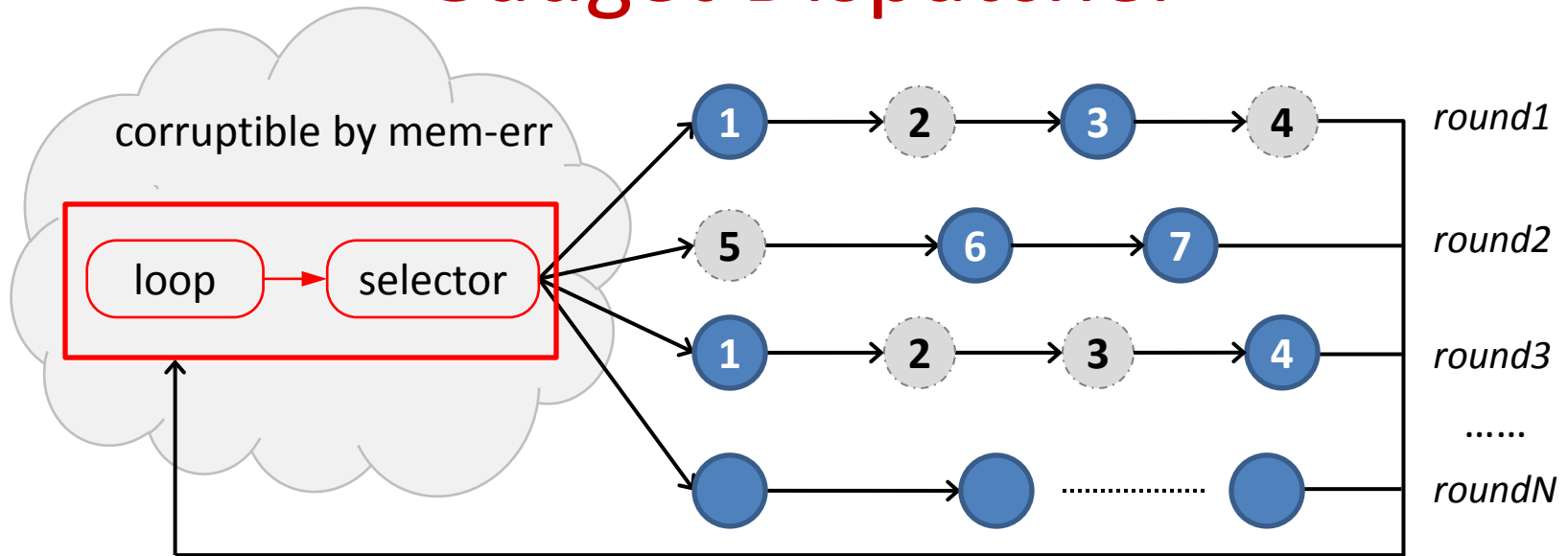
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Gadget Dispatcher



- Chain data-oriented gadgets ***“legitimately”***
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```
6  while (quota-- ) {                // loop
7      readData(sockfd, buf);        // selector
8      if(*type == NONE ) break;
9      if(*type == STREAM) *size = *(srv->cur_max);
10     else{ srv->typ = *type;  srv->total += *size; }
14 }
```

Turing-completeness

- DOP emulates a minimal language *MINDOP*
 - *MINDOP* is Turing-complete

Semantics	Statements In C	Data-Oriented Gadgets in DOP
arithmetic / logical	<code>a op b</code>	<code>*p op *q</code>
assignment	<code>a = b</code>	<code>*p = *q</code>
load	<code>a = *b</code>	<code>*p = **q</code>
store	<code>*a = b</code>	<code>**p = *q</code>
jump	<code>goto L</code>	<code>vpc = &input</code>
conditional jump	<code>if (a) goto L</code>	<code>vpc = &input if *p</code>
<code>p – &a; q – &b; op – any arithmetic / logical operation</code>		

Attack Construction

```
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- Gadget identification
 - statically identify load-semantics-store chain from LLVM IR

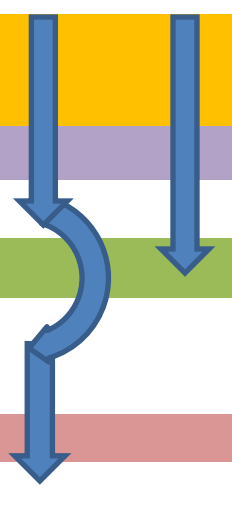
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 - static identify loops with gadgets from LLVM IR

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```



The diagram illustrates control flow within the while loop. Two vertical blue arrows point downwards from the top of the loop body. One arrow is straight and points to the 'break;' statement on line 8. The other arrow is curved and points to the 'else {' block on line 11. This represents the two possible execution paths from the loop's condition check.

- Gadget identification
 - statically identify load-semantics-store chain from LLVM IR
- Dispatcher identification
 - static identify loops with gadgets from LLVM IR
- Gadget stitching
 - select gadgets and dispatchers (manual)
 - check stitchability (manual)

Evaluation

Evaluation – Feasibility

9 x86 programs with 9 vulnerabilities

- Nginx, ProFTPD, Wu-FTPd, sshd, bitcoind,
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- 1443 in total, 110 reachable from selected CVEs

- 2 programs can build Turing-complete attack

- 3 end-to-end attacks

Case Study: Bypassing Randomization

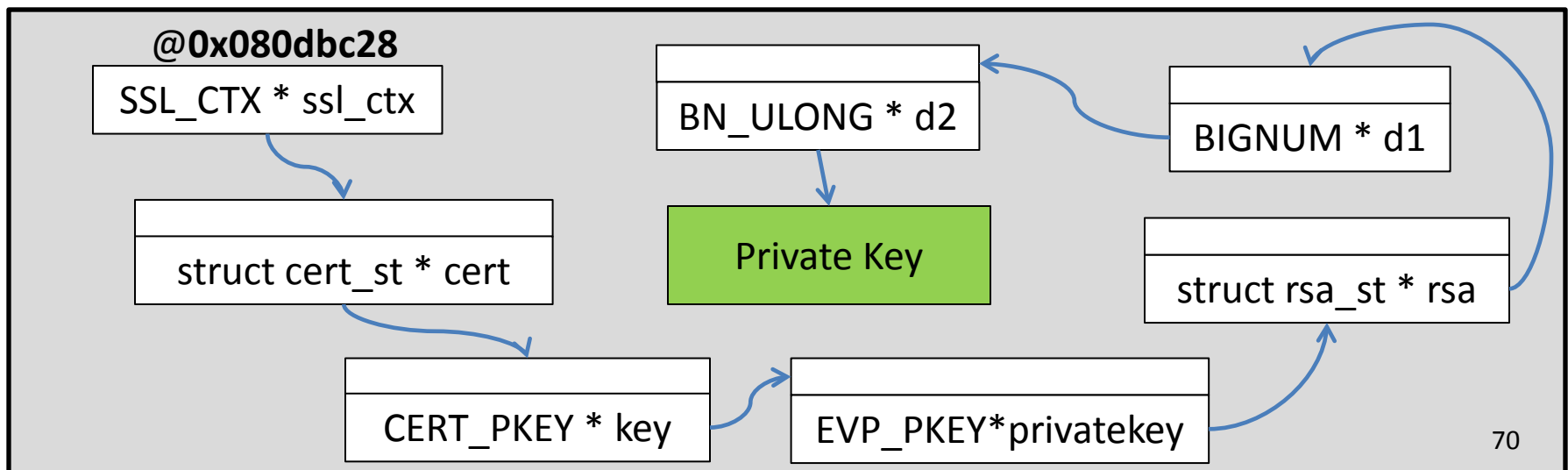
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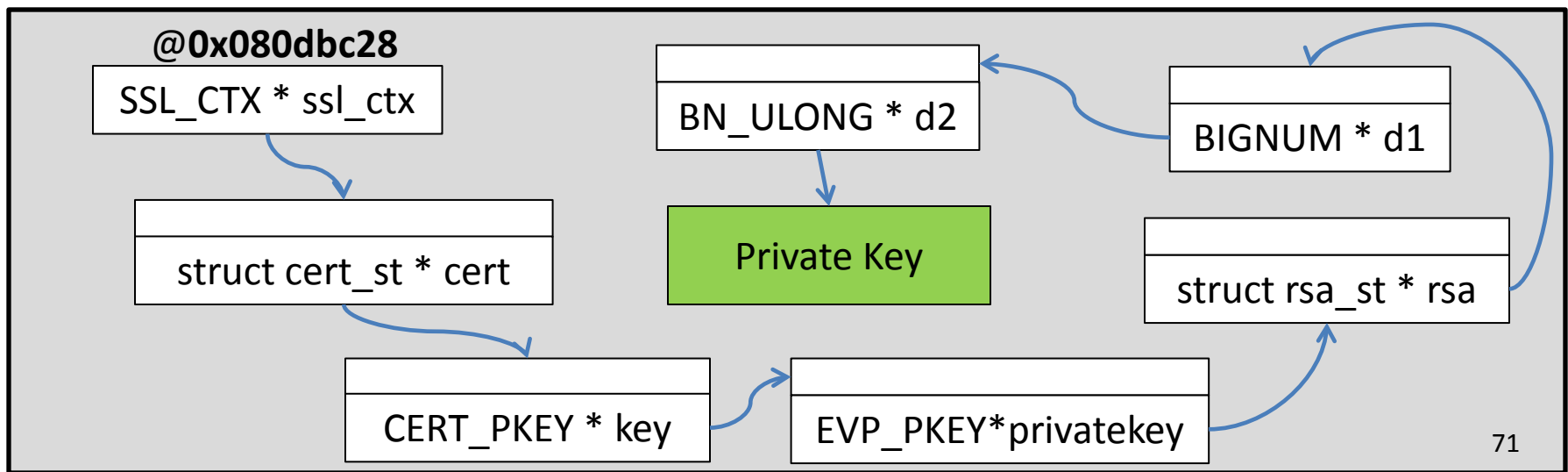
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- Gadgets

MOV	$*p = *q$
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Case Study: Bypassing Randomization

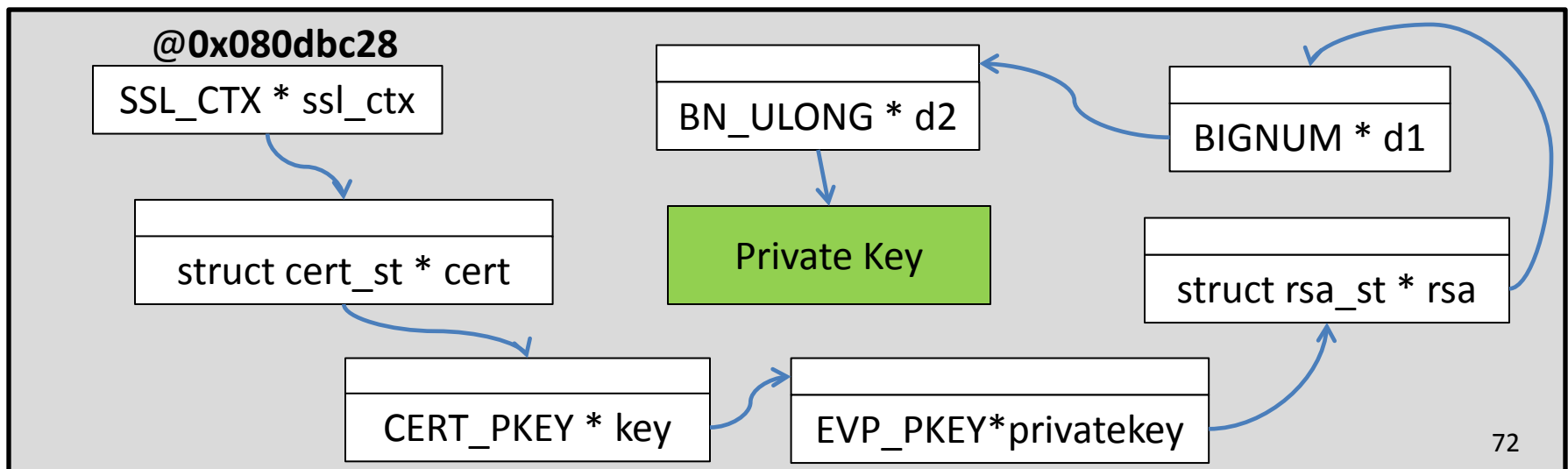
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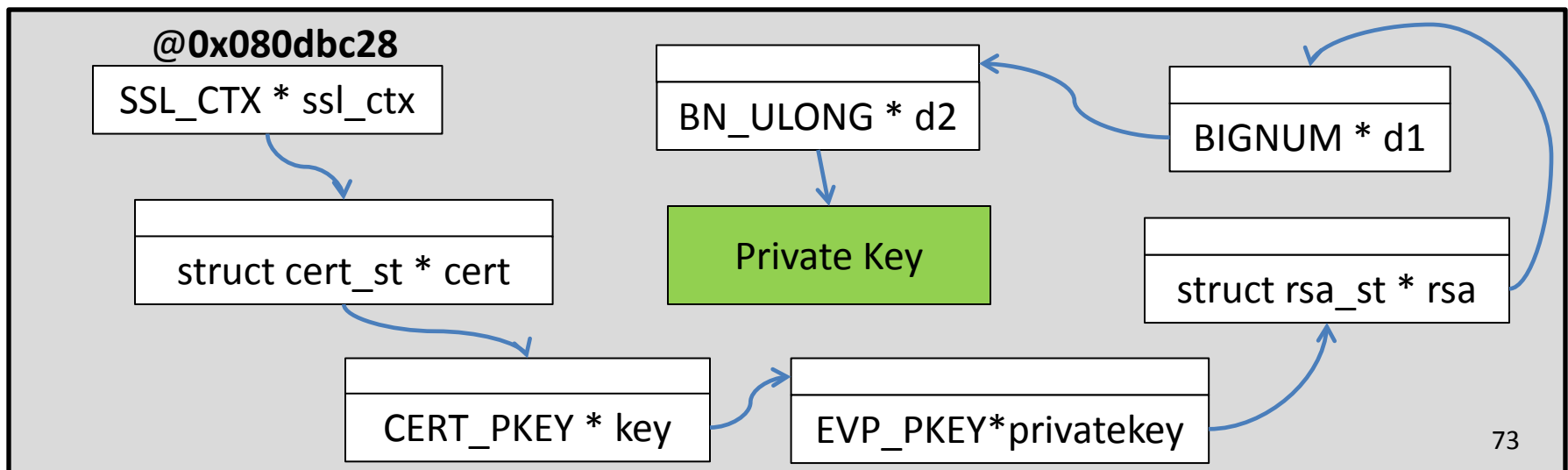
```
while (1) {  
    user_request =  
        get_user_request();  
    dispatch(user_request);  
}
```

func1() { memory_error; MOV; }
func2() { ADD; }
func3() { LOAD; }



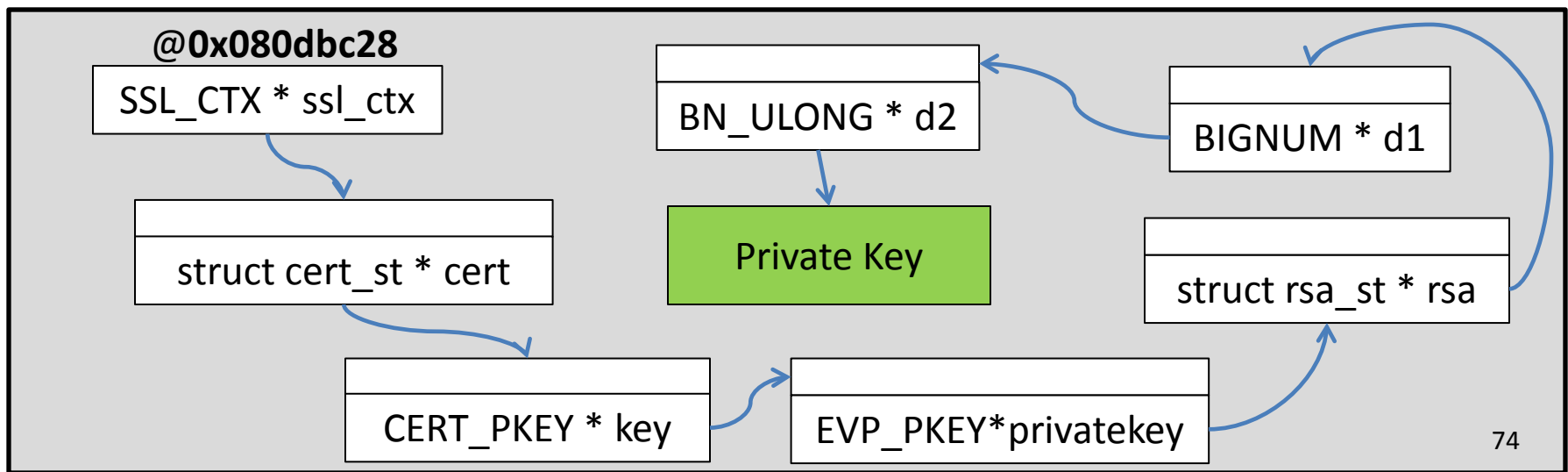
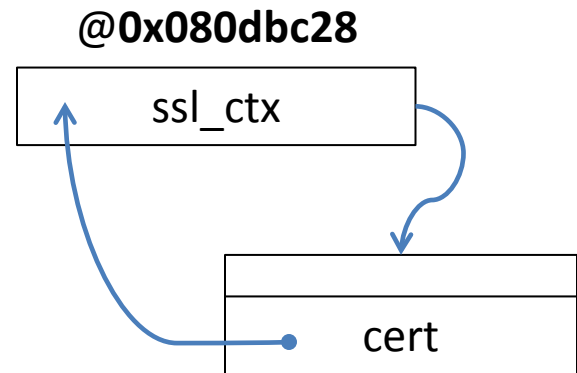
Case Study: Bypassing Randomization

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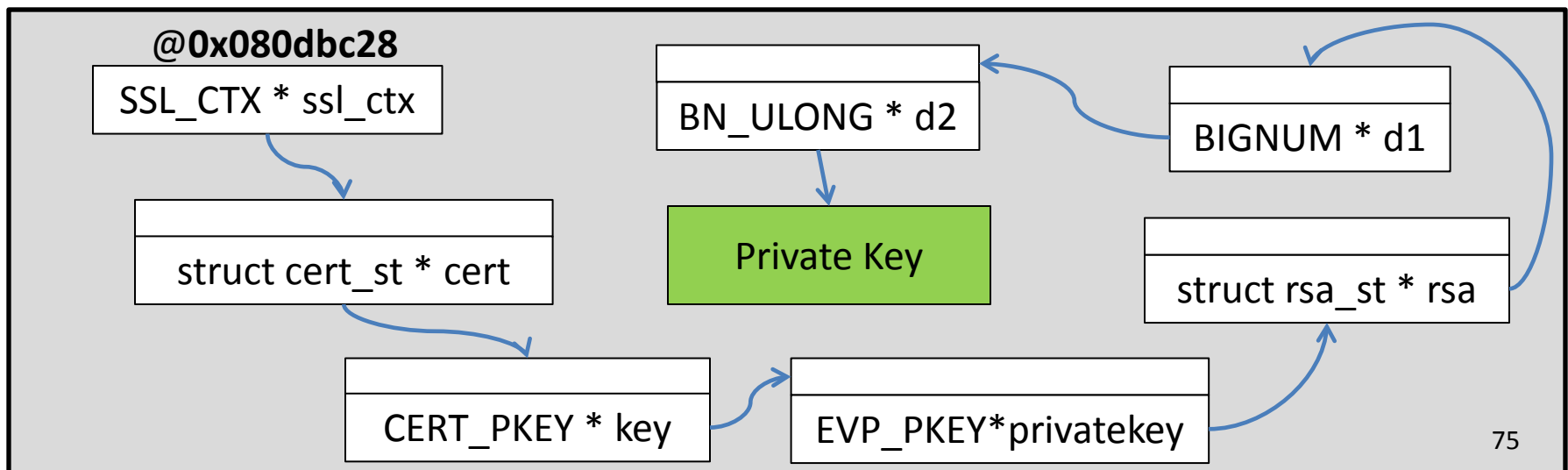
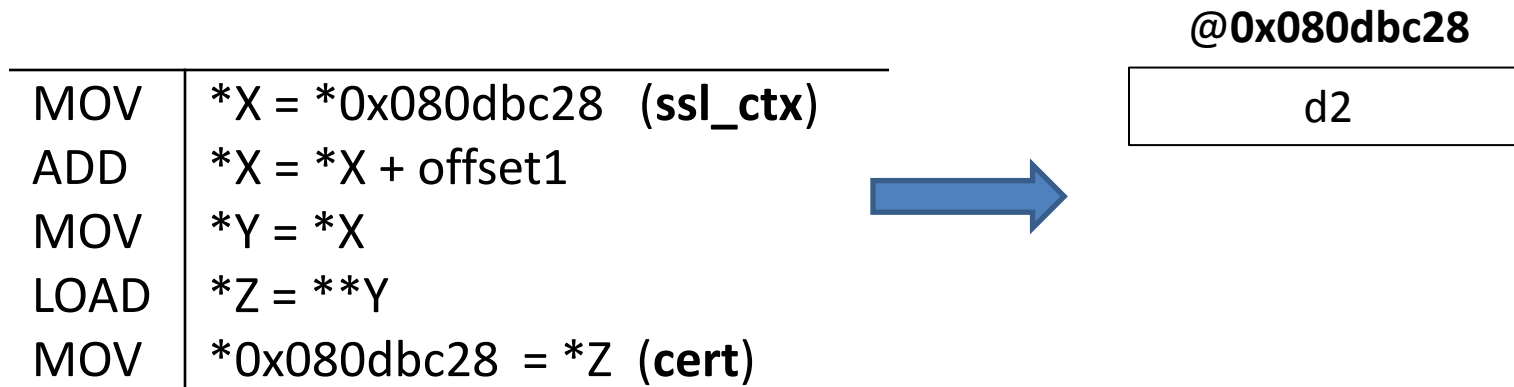


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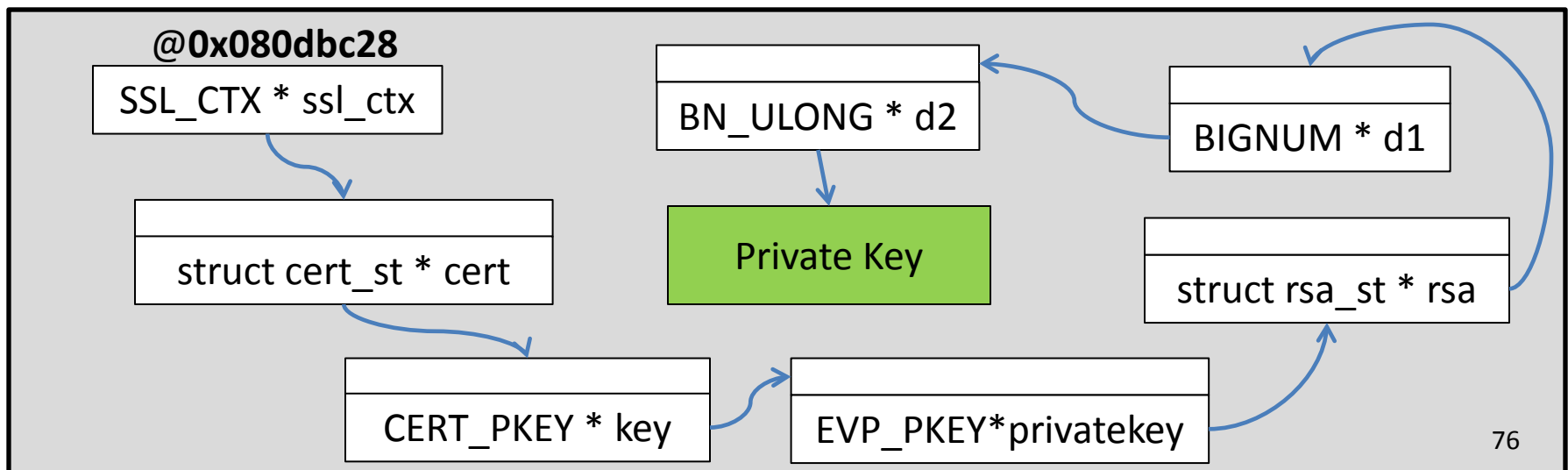
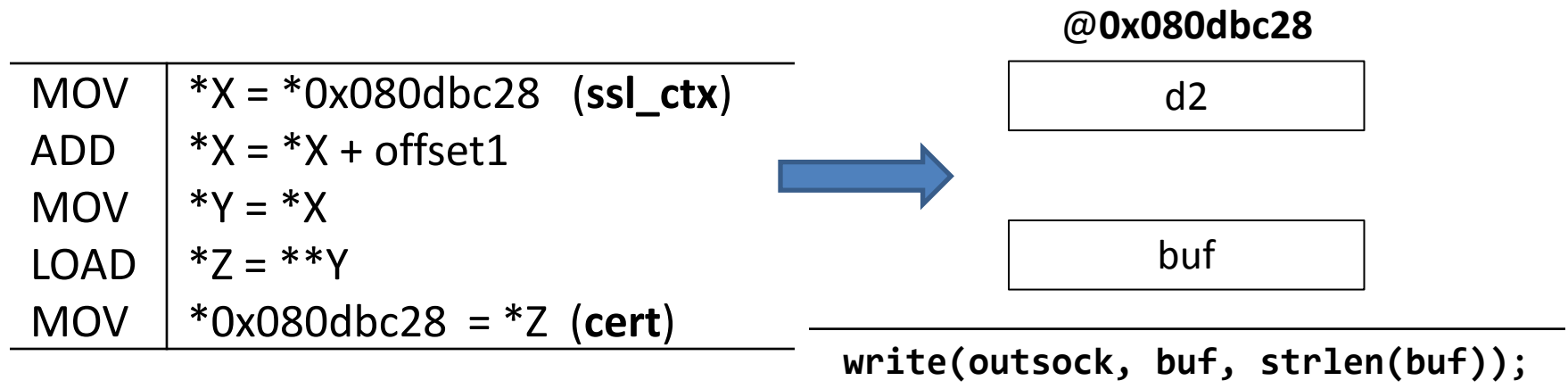
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Case Study: Bypassing Randomization

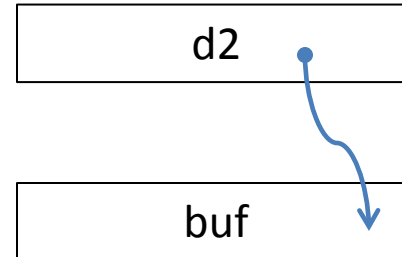


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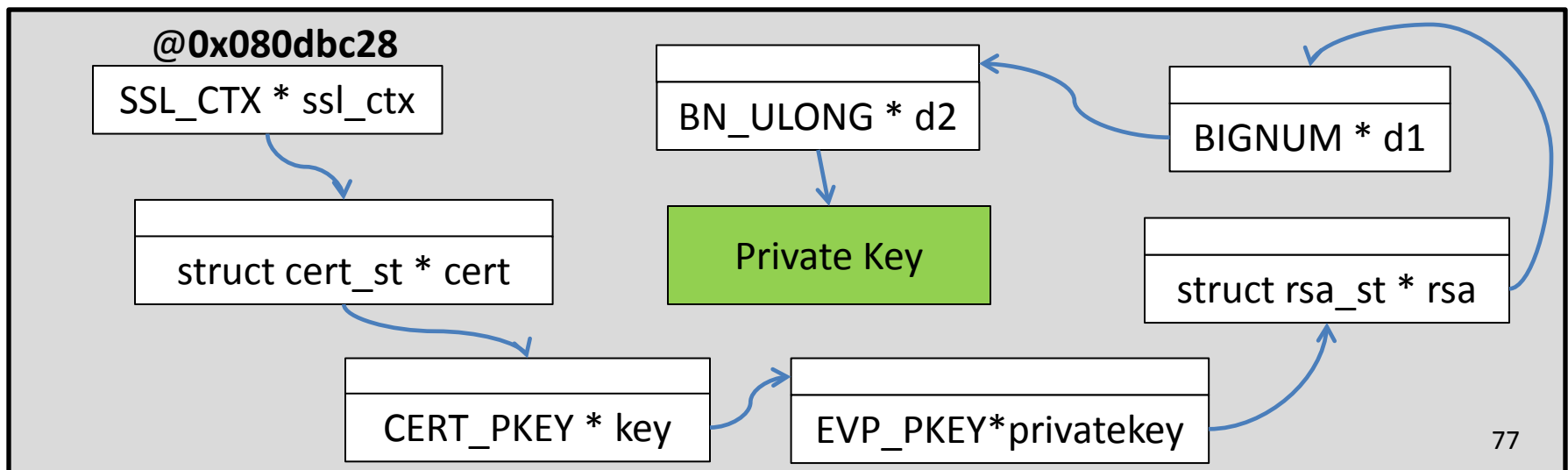
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@0x080dbc28



write(outsock, buf, strlen(buf));

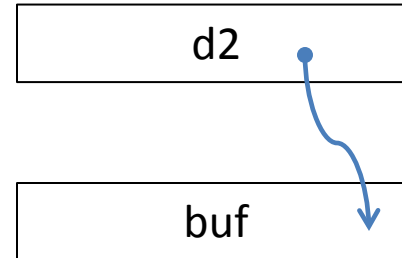


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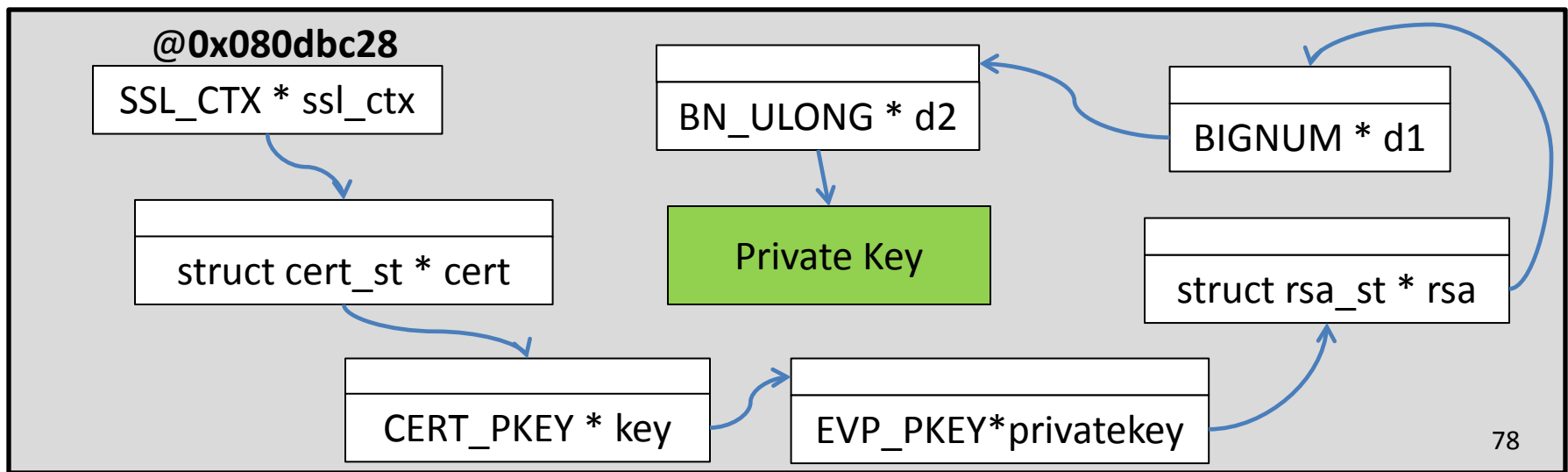


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leak private key to network

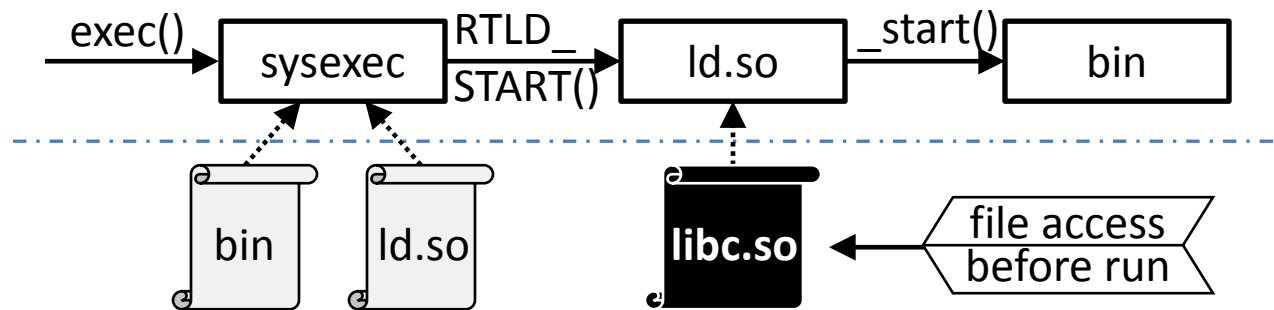


dlopen() – Dynamic Linking Interface

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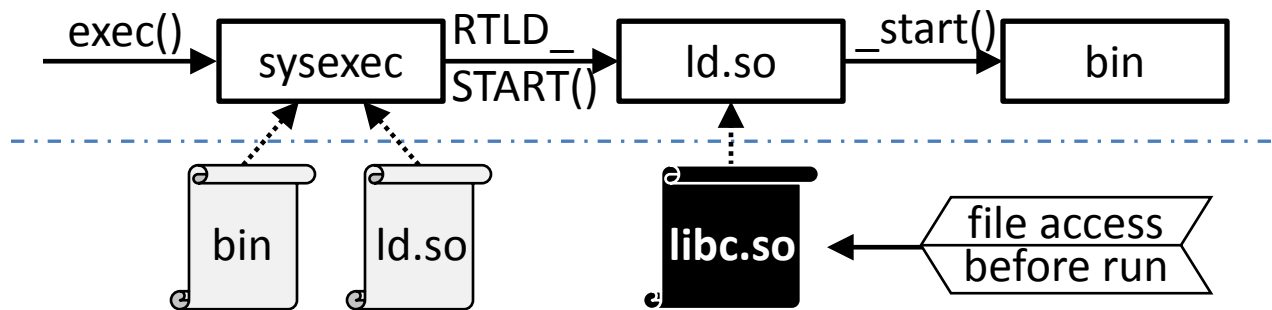
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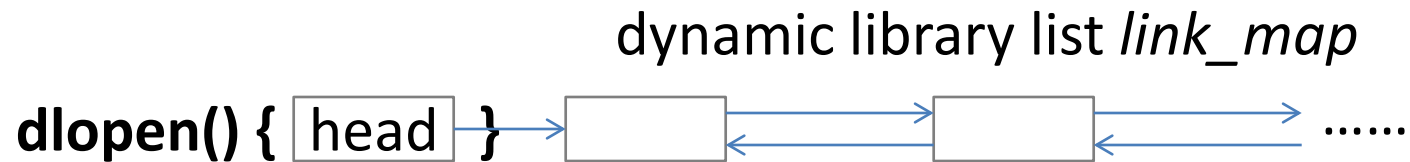


- The same to *dlopen()*

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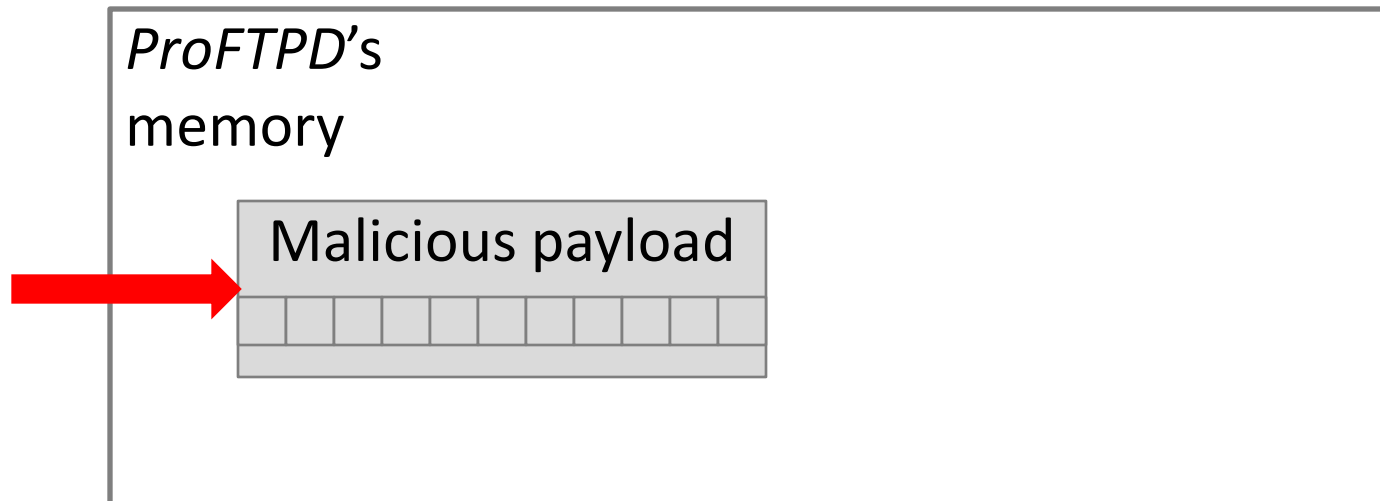
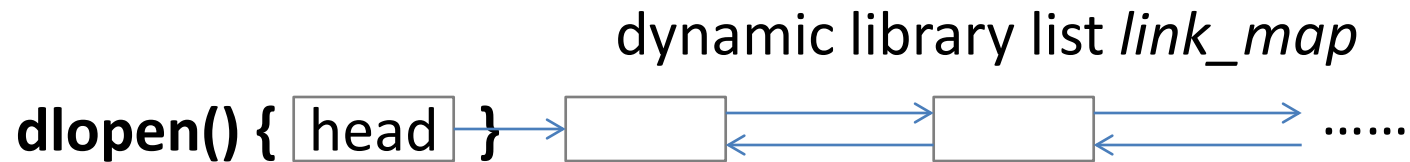
Case Study: Simulating A Network Bot

- Attacks with *dlopen*



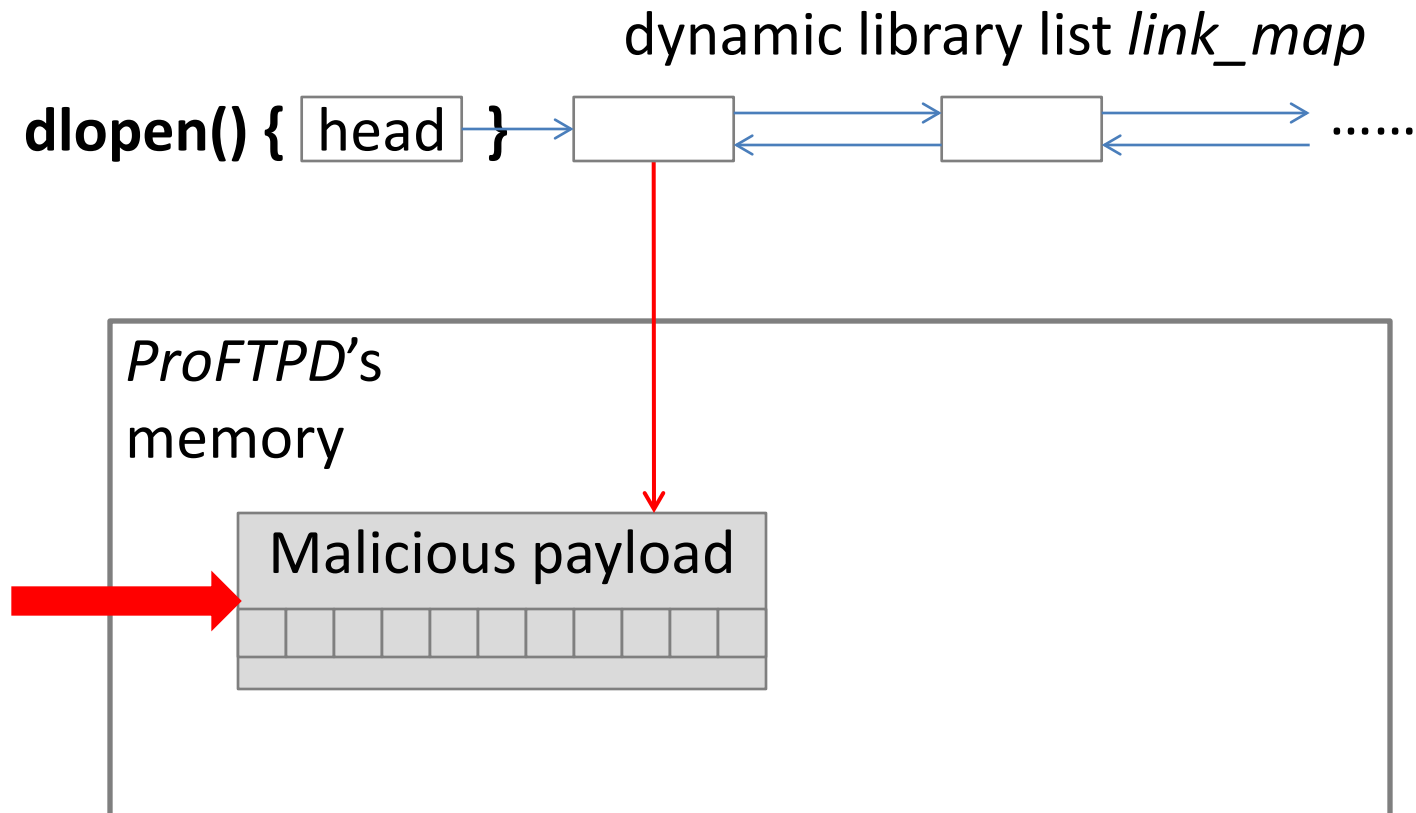
Case Study: Simulating A Network Bot

- Attacks with *dlopen*
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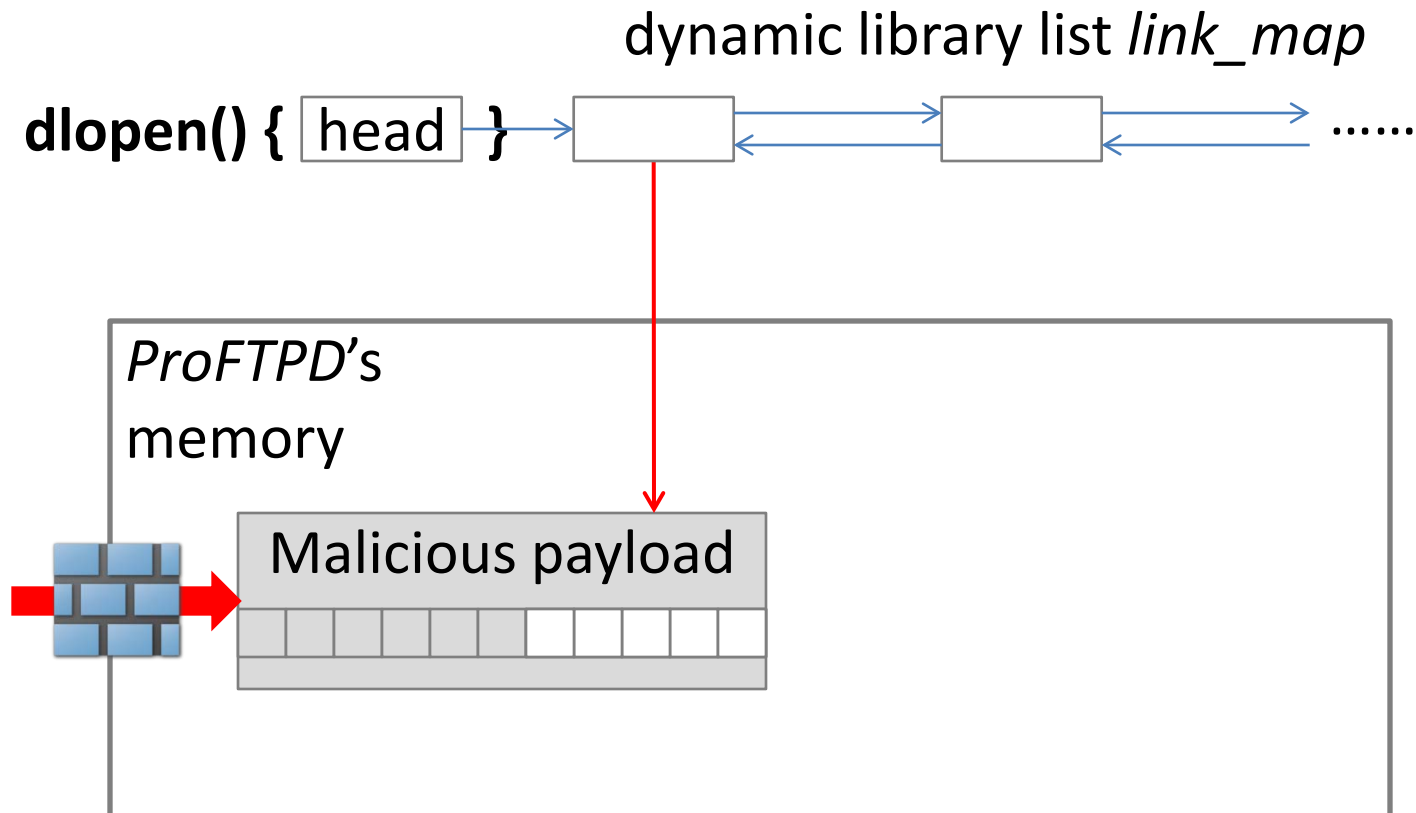
Case Study: Simulating A Network Bot

- Attacks with *dlopen*
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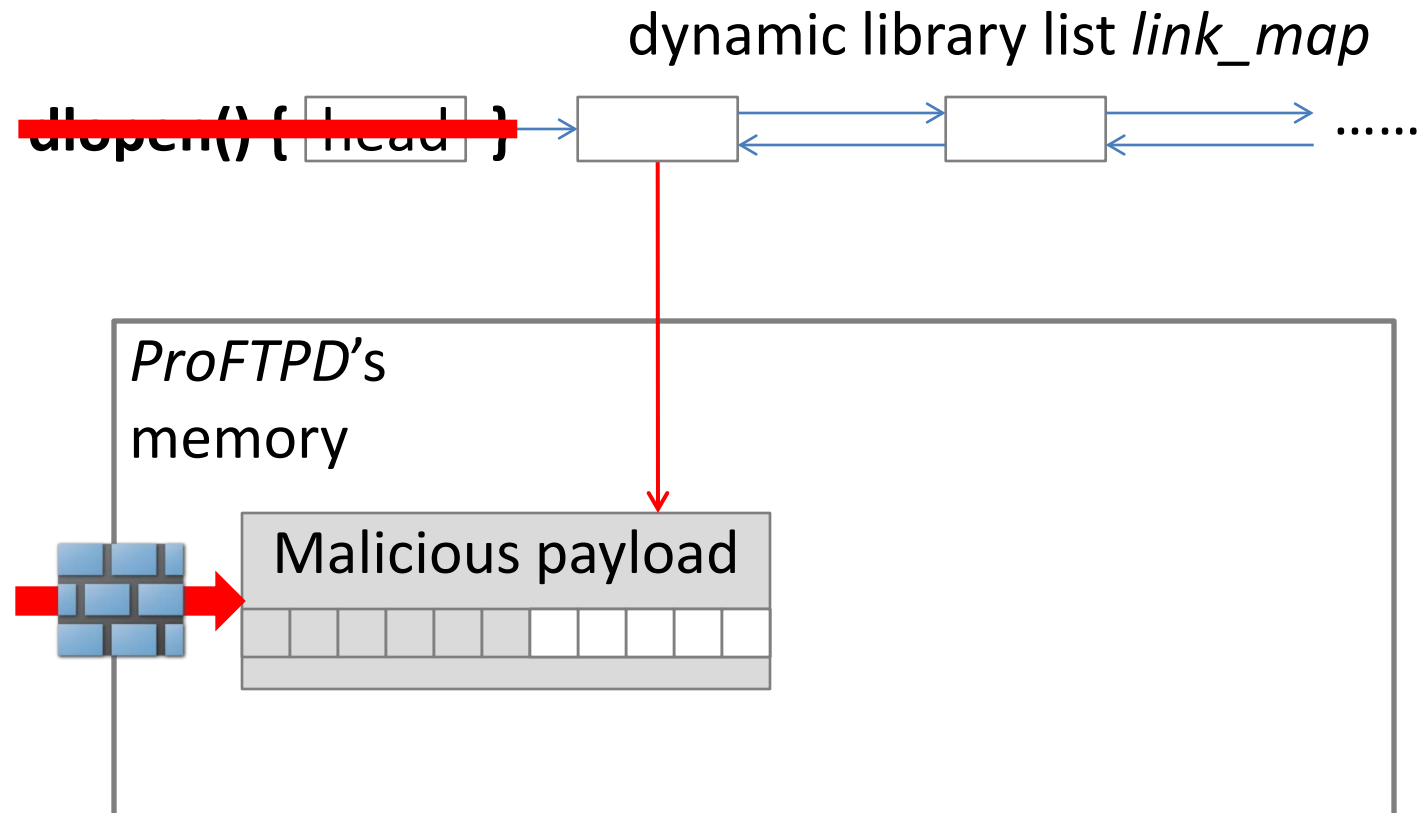
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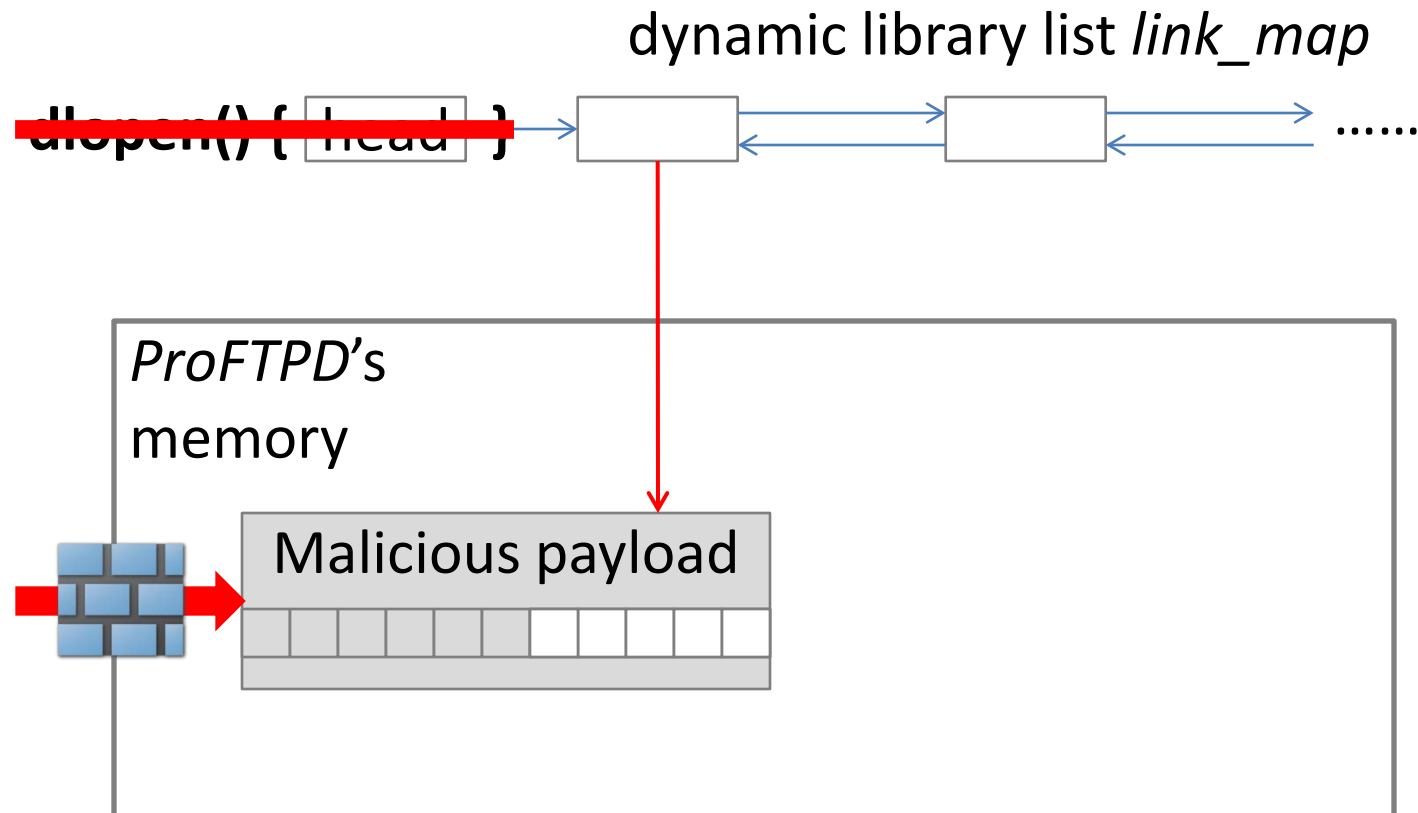
Case Study: Simulating A Network Bot

- Attacks with *dlopen*
 - send malicious payload invalid input
 - corrupt link list & call *dlopen* no call to *dlopen*



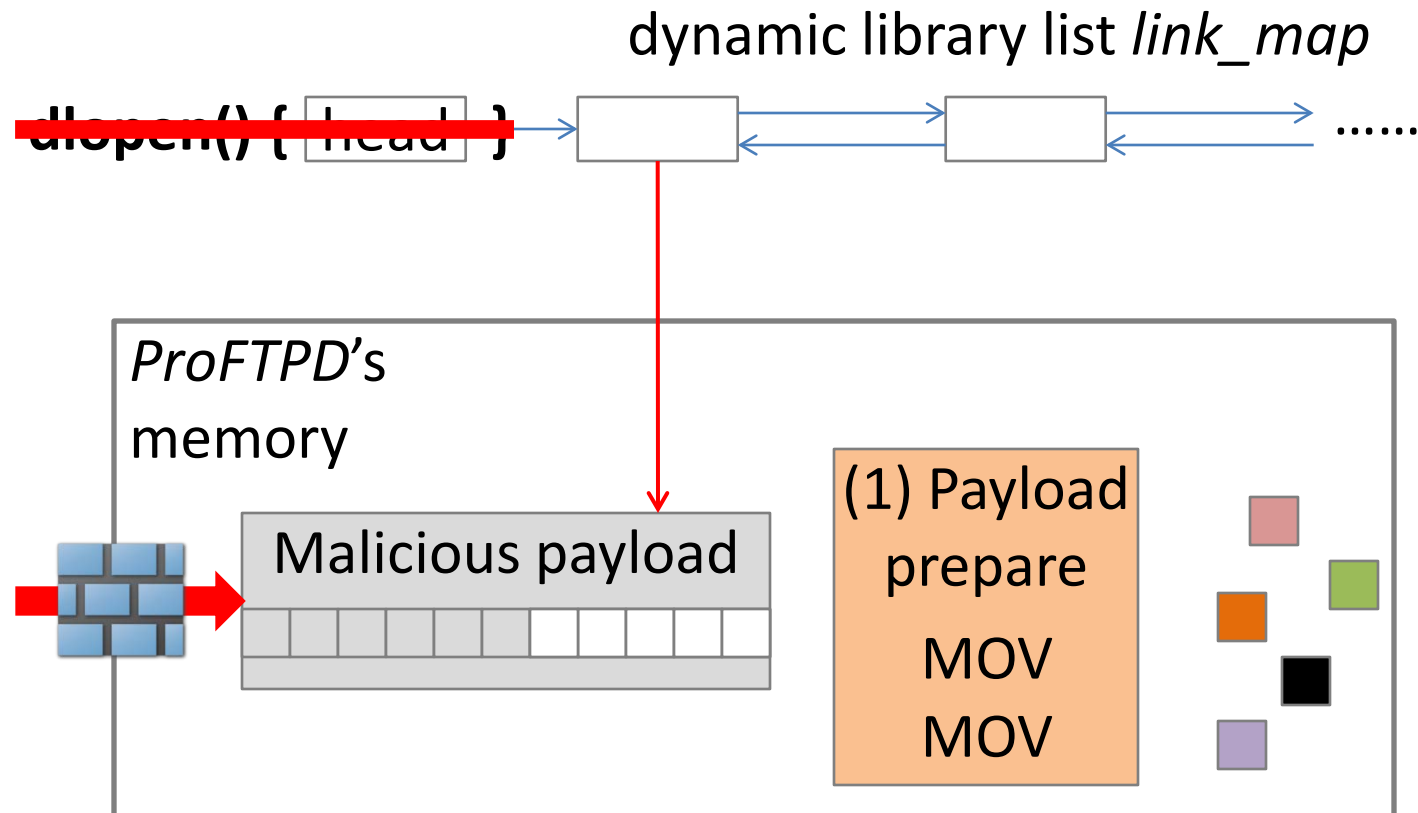
Case Study: Simulating A Network Bot

- DOP attack addresses the problems
 - send malicious payload **invalid input**
 - corrupt link list & call *dlopen* **no call to *dlopen***



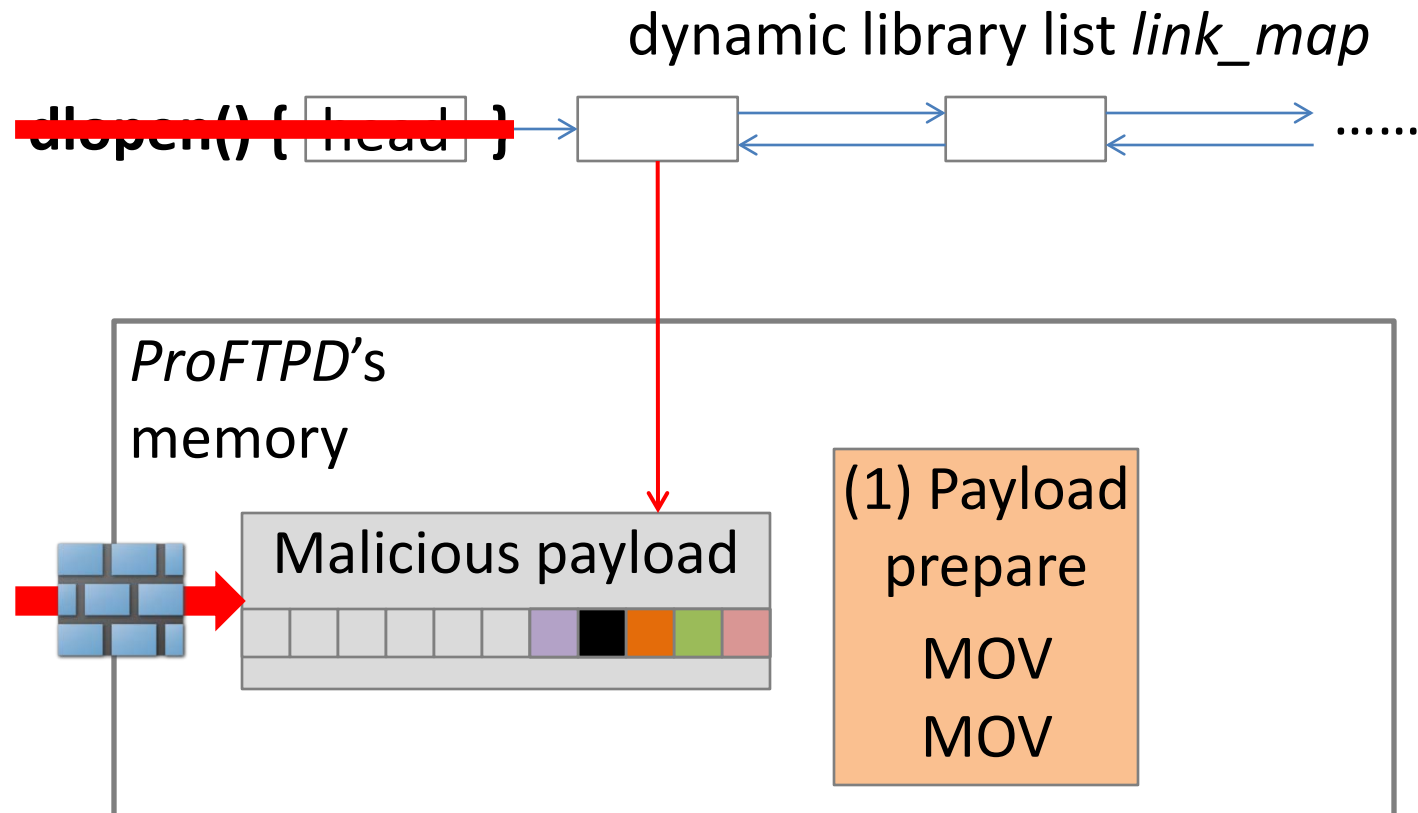
Case Study: Simulating A Network Bot

- DOP attack addresses the problems
 - construct payload in memory **invalid input**
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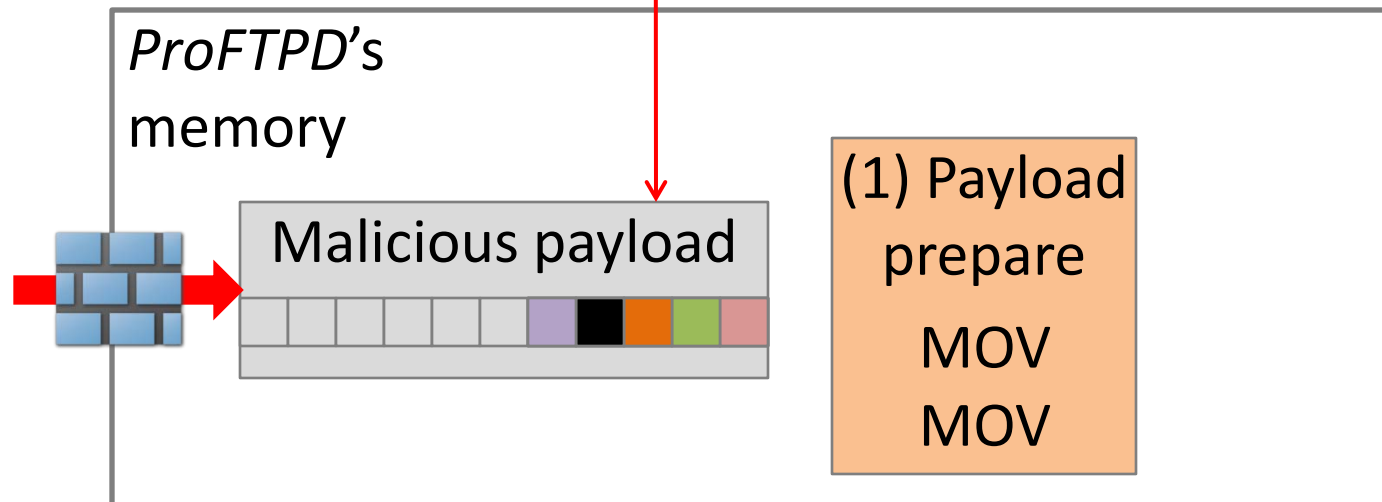


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 - force call to *dlopen* no call to *dlopen*

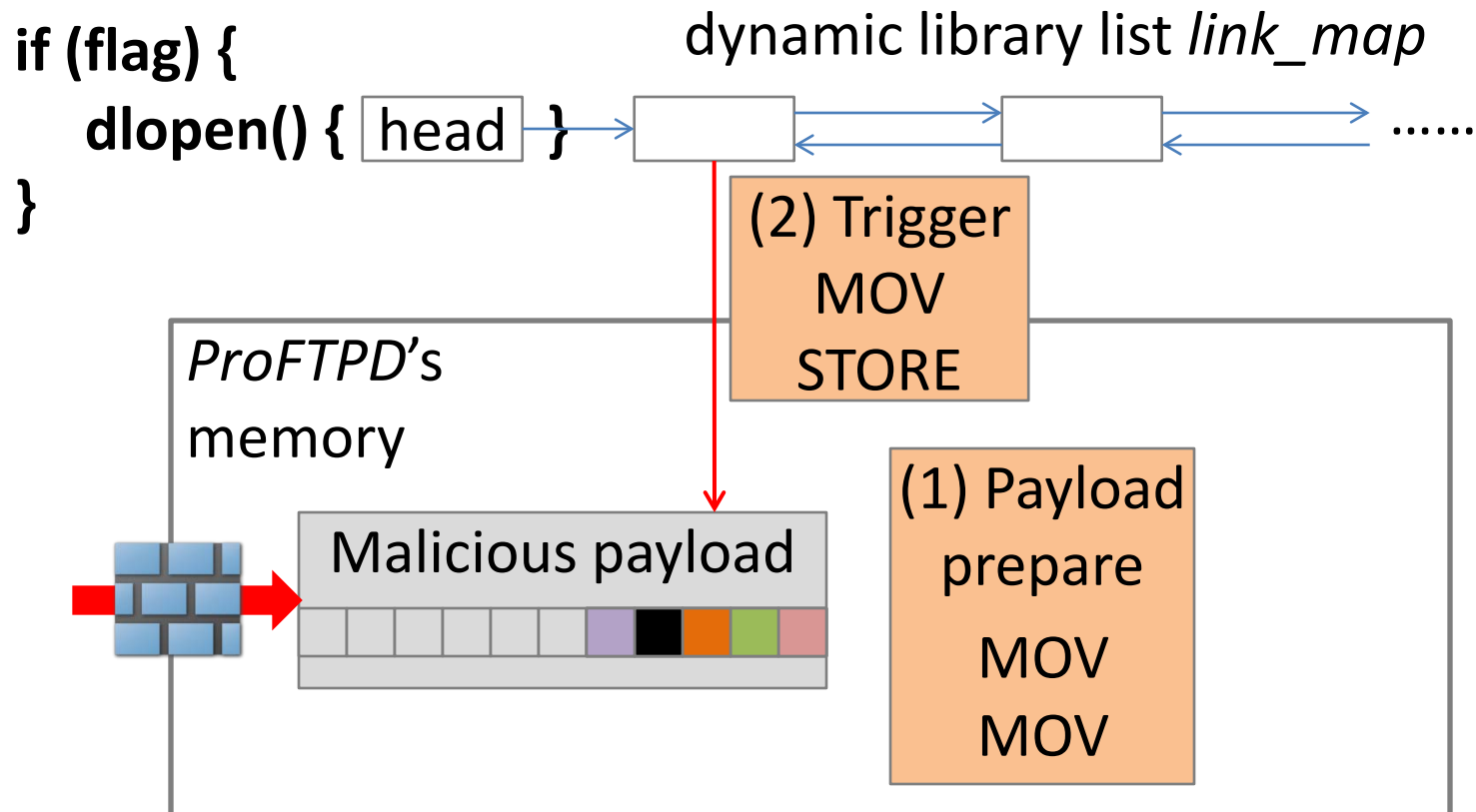
```
if (flag) {
    dlopen() { head } → [ ] ↔ [ ] ↔ .....
}
```

dynamic library list *link_map*



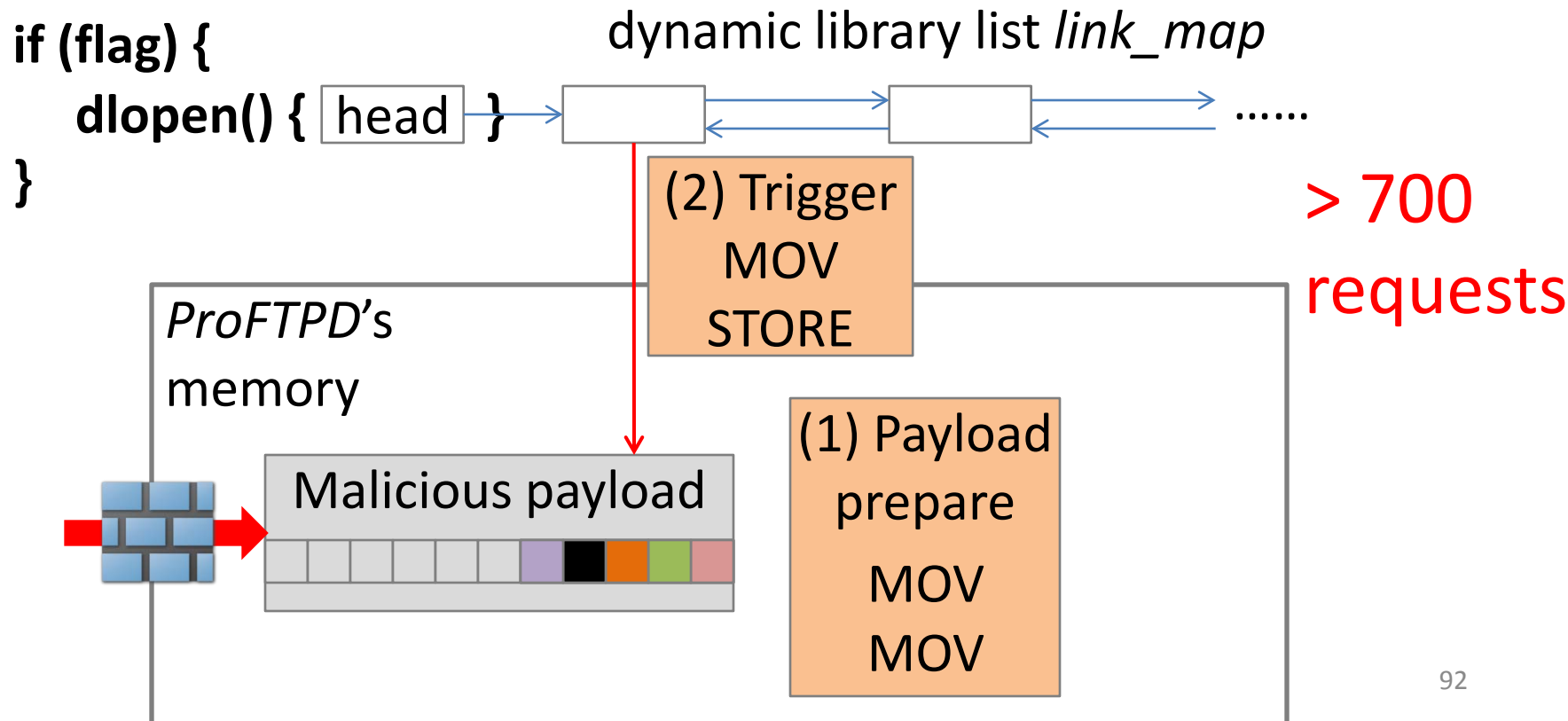
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 - *dlopen*(code_addr, shellcode)
- Code injection is back!

Related Work

Techniques	Turing Complete?	Preserve CFI?	Independent of specific data / funcs?
Non-control Data Attacks (Chen <i>et al.</i> 2005)		✓	
COOP (Schuster <i>et al.</i> 2015)	✓		✓
FlowStitch (Hu <i>et al.</i> 2015)		✓	
Printf-Oriented Programming (Carlini <i>et al.</i> 2015)	✓	✓	
Control Jujustu (Evans <i>et al.</i> 2015)		✓	
Data-Oriented Programming	✓	✓	✓

Potential Defenses

- Memory Safety
 - e.g., Cyclone (Jim *et al.* 2002), CCured (Necula *et al.* 2002) , SoftBounds+CETS (Nagarakatte *et al.* 2009, 2010)
 - high performance overhead (> 100%)
- Data-flow Integrity
 - e.g, DFI (Castro *et al.* 2006) , kernel DFI (Song *et al.* 2016)
- Fined-grained randomization in data space
 - e.g., DSR (Bhatkar *et al.* 2008)
- Hardware & software fault isolation
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No practical defenses yet !

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- Data-Oriented Programming (DOP)
 - build expressive non-control data attacks
 - independent of specific data / functions
- In real-world programs, DOP can build attacks
 - bypass ASLR w/o address leakage
 - simulate a network bot
 - enable code injection

Thanks!

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Non-control data attacks are available

<http://huhong-nus.github.io/advanced-DOP/>